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Artin glueings of toposes as adjoint split extensions $\stackrel{\diamond}{\approx}$



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ABSTRACT

Artin glueings of frames correspond to adjoint split extensions in the category of frames and finite-meet-preserving maps. We extend these ideas to the setting of toposes and show that Artin glueings of toposes correspond to a 2-categorical notion of adjoint split extension in the 2-category of toposes, finite-limit-preserving functors and natural transformations. A notion of morphism between these split extensions is introduced, which allows the category $\text{Ext}(\mathcal{H}, \mathcal{N})$ to be constructed. We show that $\text{Ext}(\mathcal{H}, \mathcal{N})$ is equivalent to $\text{Hom}(\mathcal{H}, \mathcal{N})^{\text{op}}$, and moreover, that this can be extended to a 2-natural contravariant equivalence between the Hom 2-functor and a naturally defined Ext 2-functor.

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1. Introduction

Artin glueings of toposes were introduced in [1] and provide a way to view a topos \mathcal{G} as a combination of an open subtopos $\mathcal{G}_{\mathfrak{o}(U)}$ and its closed complement $\mathcal{G}_{\mathfrak{c}(U)}$. This situation may be described as the 'internal' view, but we might instead look at it externally. Here we have that Artin glueings of two toposes \mathcal{H} and \mathcal{N} correspond to solutions to the problem of which toposes \mathcal{G} does \mathcal{H} embed in as an open subtopos and \mathcal{N} as its closed complement.

There is an analogy to be made with semidirect products of groups. We may either view a group as being generated in a natural way from two complemented subgroups (one of which is normal), or externally, view a semidirect product as a solution to the problem of how to embed groups H and N as complemented

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subobjects so that N is normal. Of particular importance to us is that semidirect products correspond to split extensions of groups (up to isomorphism).

Artin glueings of Grothendieck toposes decategorify to the setting of frames and in this algebraic setting the analogy to semidirect products has been made precise. In [11] it was shown that Artin glueings correspond to certain split extensions in the category of frames with finite-limit-preserving maps. While these results were proved in the setting of frames, it is not hard to see that the arguments carry over to Heyting algebras. It is this view that we now extend back to the elementary topos setting. We now recall the main results of [11].

In the category of frames with finite-meet-preserving maps, there exist zero morphisms given by the constant 'top' maps. This allows us to consider kernels and cokernels. Cokernels always exist and the cokernel of $f: N \to G$ is given by $e: G \to \downarrow f(0)$ where $e(g) = f(0) \land g$. This map has a right adjoint splitting e_* sending h to $h^{f(0)}$. Kernels do not always exist, but kernels of cokernels always do, and the kernel of $e: G \to \downarrow u$ is the inclusion of $\uparrow u \subseteq G$. The cokernel is readily seen to be the open sublocale corresponding to u and the kernel the corresponding closed sublocale. This immediately gives that the split extensions whose splittings are adjoint to the cokernel correspond to Artin glueings.

A notion of protomodularity for categories equipped with a distinguished class of split extensions was first introduced in [5]. In a similar way, in [11] (as well as the current paper) a distinguished role is played by those split extensions whose splittings are right adjoint to the cokernel.

With this correspondence established, the corresponding Ext functor was shown to be naturally isomorphic to the Hom functor. Each hom-set Hom(H, N) has an order structure and this order structure was shown to correspond contravariantly in Ext(H, N) to morphisms of split extensions. Finally, it was demonstrated how the meet operation in Hom(H, N) naturally induces a kind of 'Baer sum' in Ext(H, N).

In this paper all of the above results find natural generalisation to the topos setting after we provide definitions for the analogous 2-categorical concepts. (Also see the paper [16] by Niefield which appeared after we wrote this paper and discusses the relationship between (non-split) extensions and glueing in a very general context. The current paper uses split extensions and discusses the functorial nature of the construction for toposes in more detail.)

2. Background

2.1. 2-categorical preliminaries

There are a number of 2-categories and 2-categorical constructions considered in this paper and so we provide a brief overview of these here. All the 2-categories we consider shall be strict.

Briefly, a (strict) 2-category consists of objects, 1-morphisms between objects and 2-morphisms between 1-morphisms. Phrased another way, instead of hom-sets between objects as is the case with 1-categories, for any two objects A and B we have an associated hom-category Hom(A, B).

Both 1-morphisms and 2-morphisms may be composed under the right conditions. If $F: A \to B$ and $G: B \to C$ are 1-morphisms, then we may compose them to yield $GF: A \to C$. We usually represent this with juxtaposition, though if an expression is particularly complicated we may use $G \circ F$. As with natural transformations, there are two ways to compose 2-morphisms — vertically and horizontally. We may compose $\alpha: F \to G$ and $\beta: G \to H$ vertically to give $\beta \alpha: F \to H$, sometimes written $\beta \circ \alpha$. Orthogonally, if we have $F_2F_1: A \to C, G_2G_1: A \to C, \alpha: F_1 \to G_1$ and $\beta: F_2 \to G_2$, then we may compose α and β horizontally to form $\beta * \alpha: F_2F_1 \to G_2G_1$. Vertical and horizontal composition are related by the so-called interchange law. For each object A there exists an identity 1-morphism id_A and for each 1-morphism F there exists an identity 2-morphism id_F.

Just as one can reverse the arrows of a category \mathcal{B} to give \mathcal{B}^{op} , one can reverse the 1-morphisms of a 2-category \mathcal{C} to give \mathcal{C}^{op} . It is also possible to reverse the directions of the 2-morphisms yielding \mathcal{C}^{co} and when both the 1-morphisms and 2-morphisms are reversed we obtain $\mathcal{C}^{\text{co op}}$.

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In this paper we will make extensive use of string diagrams. For an introduction to string diagrams for 2-categories see [14]. We will use the convention that vertical composition is read from bottom to top and horizontal composition runs diagrammatically from left to right.

We consider 2-functors between 2-categories defined as follows. (We follow the convention that 2-functors are not necessarily strict.)

Definition 2.1. A 2-functor \mathcal{F} between 2-categories \mathcal{C} and \mathcal{D} consists of a function \mathcal{F} sending objects Cin \mathcal{C} to objects $\mathcal{F}(C)$ in \mathcal{D} and for each pair of objects C_1 and C_2 in \mathcal{C} a functor \mathcal{F}_{C_1,C_2} : Hom $(C_1,C_2) \to$ Hom $(\mathcal{F}(C_1),\mathcal{F}(C_2))$, for which we use the same name. Additionally, for each pair of composable 1-morphisms (F,G) we have an invertible 2-morphism $\omega_{G,F} \colon \mathcal{F}(G) \circ \mathcal{F}(F) \to \mathcal{F}(G \circ F)$ called the compositor, and for each object C in \mathcal{C} we have an invertible 2-morphism $\kappa_A \colon \mathrm{Id}_{\mathcal{F}(A)} \to \mathcal{F}(\mathrm{Id}_A)$ called the unitor. This data satisfies the following constraints.

- (1) Let $\alpha: F_1 \to F_2$ and $\beta: G_1 \to G_2$ be horizontally composable 2-morphisms. Then the compositors must satisfy the naturality condition $\omega_{G_2,F_2}(\mathcal{F}(\beta) * \mathcal{F}(\alpha)) = \mathcal{F}(\beta * \alpha) \omega_{G_1,F_1}$.
- (2) The compositors must be associative in the sense that if $F: X \to Y$, $G: Y \to Z$ and $H: Z \to W$ are 1-morphisms, then $\omega_{H,GF}(\operatorname{id}_{\mathcal{F}(H)} * \omega_{G,F}) = \omega_{HG,F}(\omega_{H,G} * \operatorname{id}_{\mathcal{F}(F)}).$
- (3) If $F: X \to Y$ is a 1-morphism, then we have the unit axiom $\omega_{F, \mathrm{Id}_X}(\mathrm{id}_{\mathcal{F}(F)} * \kappa_X) = \mathrm{id}_{\mathcal{F}(F)} = \omega_{\mathrm{Id}_Y, F}(\kappa_Y * \mathrm{id}_{\mathcal{F}(F)}).$

There is a notion of 2-natural transformation between 2-functors defined as follows.

Definition 2.2. Let $(\mathcal{F}_1, \omega_1, \kappa_1), (\mathcal{F}_2, \omega_2, \kappa_2) \colon \mathcal{X} \to \mathcal{Y}$ be two 2-functors. A 2-natural transformation $\rho \colon \mathcal{F}_1 \to \mathcal{F}_2$ is given by two families:

- (1) A 1-morphism $\rho_X \colon \mathcal{F}_1(X) \to \mathcal{F}_2(X)$ for each object X in \mathcal{X} .
- (2) An invertible 2-morphism $\rho_F \colon \mathcal{F}_2(F)\rho_X \to \rho_Y \mathcal{F}_1(F)$ for each 1-morphism $F \colon X \to Y$ in \mathcal{X} .

They must satisfy the following coherence conditions. First if $F: X \to Y$ and $G: Y \to Z$ are 1-morphisms in \mathcal{X} , then ρ must respect composition, so that the following diagram commutes.



Next for each object $X \in \mathcal{X}$, ρ must respect the identity Id_X . For this we need the following diagram to commute.



Finally, they must satisfy the following 'naturality' condition for $F, F': X \to Y$ and $\alpha: F \to F'$.

$$\begin{array}{c|c} \mathcal{F}_{2}(F) \circ \rho_{X} & \xrightarrow{\rho_{F}} \rho_{Y} \circ \mathcal{F}_{1}(F) \\ \\ \mathcal{F}_{2}(\alpha)\rho_{X} & & & \\ & & & \\ \mathcal{F}_{2}(F') \circ \rho_{X} & \xrightarrow{\rho_{F'}} \rho_{Y} \circ \mathcal{F}_{1}(F') \end{array}$$

When each component ρ_X is an equivalence, we call ρ a 2-natural equivalence.

One 2-functor of note is the 2-functor Op: $\operatorname{Cat}^{\operatorname{co}} \to \operatorname{Cat}$ which sends a category \mathcal{C} to its opposite category $\mathcal{C}^{\operatorname{op}}$. We will use this 2-functor in Section 5 to help compare two 2-functors of different variances.

Limits and colimits have 2-categorical analogues, which will be used extensively throughout this paper. A more complete introduction to these concepts can be found in [13]. In particular, we will make use of 2-pullbacks and 2-pushouts, as well as comma and cocomma objects, which we describe concretely below.

Definition 2.3. Given two 1-morphisms $F: \mathcal{B} \to \mathcal{D}$ and $G: \mathcal{C} \to \mathcal{D}$ their *comma object* is shown in the following diagram



and satisfies the following two conditions.

(1) Let $T: \mathcal{X} \to \mathcal{B}$ and $S: \mathcal{X} \to \mathcal{C}$ be 1-morphisms and let $\psi: FT \to GS$ be a 2-morphism. Then there exists a 1-morphism $H: \mathcal{X} \to \mathcal{P}$ and invertible 2-morphisms $\nu: P_GH \to T$ and $\mu: P_FH \to S$ satisfying that $G\mu \circ \varphi H \circ F\nu^{-1} = \psi$.



(2) If $H, K: \mathcal{X} \to \mathcal{P}$ are 1-morphisms and $\alpha: P_G H \to P_G K$ and $\beta: P_F H \to P_F K$ are 2-morphisms satisfying that $\varphi K \circ F \alpha = G \beta \circ \varphi H$, then there exists a unique 2-morphism $\gamma: H \to K$ such that $P_G \gamma = \alpha$ and $P_F \gamma = \beta$.

A 2-pullback is defined similarly, except both φ and ψ are required to be invertible, and is represented as follows.



Cocomma objects and 2-pushouts may be defined dually.

We now recall the definition of a fibration of categories.

Definition 2.4. Let $F: \mathcal{X} \to \mathcal{Y}$ be a functor. A morphism $\overline{f}: \overline{A} \to \overline{B}$ in \mathcal{X} is *cartesian* with respect to F if for any $\overline{g}: \overline{C} \to \overline{B}$ in \mathcal{X} and $h: F(\overline{C}) \to F(\overline{A})$ in \mathcal{Y} with $F(\overline{g}) = F(\overline{f})h$, there exists a unique map $\overline{h}: \overline{C} \to \overline{A}$ with $F(\overline{h}) = h$ and $\overline{fh} = \overline{g}$. We say $F: \mathcal{X} \to \mathcal{Y}$ is a *(Street) fibration* in Cat if for any morphism $f: A \to F(\overline{B})$ in \mathcal{Y} there exists a cartesian lifting $\overline{f}: \overline{A} \to \overline{B}$ and an isomorphism $j: F(\overline{A}) \cong A$ with $F(\overline{f}) = fj$.

In fact, we will also need the notion of a fibration in other 2-categories, such as the 2-category Cat_{lex} of finitely-complete categories and finite-limit-preserving functors. The general definitions of fibrations, morphisms of fibrations and 2-morphisms of fibrations can be found, for example, in [7, Definitions 3.4.3–3.4.5]. However, it is not hard to see that the fibrations in Cat_{lex} are simply the finite-limit-preserving functors which are fibrations in Cat.

2.2. Elementary toposes and Artin glueings

By a *topos* we mean an *elementary topos* — that is, a cartesian-closed category admitting finite limits and containing a subobject classifier. The usual 2-category of toposes has 1-morphisms given by *geometric morphisms* and 2-morphisms given by natural transformations.

For Grothendieck toposes, the subobjects of the terminal object may be imbued with the structure of a frame. Moreover, a geometric morphism between two Grothendieck toposes induces a locale homomorphism between their locales of subterminals. This induces a functor from the category of Grothendieck toposes into the category of locales, which is in fact a reflector.

A subtopos of a topos \mathcal{E} is a fully faithful geometric morphism $i: \mathcal{S} \hookrightarrow \mathcal{E}$. Subtoposes when acted upon by the localic reflection may sometimes be sent to open or closed sublocales. Those sent to open sublocales we call open subtoposes and those sent to closed sublocales we call closed subtoposes. Since open (or closed) sublocales correspond to elements of the frame, it follows that any open (or closed) subtopos corresponds to a particular subterminal U. This suggests a similar notion of open/closed subtopos corresponding to a particular subterminal U even in the elementary setting.

The open subtopos $\mathcal{E}_{\mathfrak{o}(U)}$ corresponding to a subterminal U has a reflector given by the exponential functor $(-)^U$. It is not hard to see that this topos is equivalent to the slice topos \mathcal{E}/U , which in turn can be thought as the full subcategory of the objects in \mathcal{E} admitting a map into U. From this point of view, the reflector $E: \mathcal{E} \to \mathcal{E}/U$ maps an object X to the product $X \times U$. We denote its right adjoint by $E_* = (-)^U$ and write θ and ε for the unit and counit respectively. Note that $E_*E(G) = (G \times U)^U \cong G^U$. In addition to a right adjoint, E also has a left adjoint E_1 , which is simply the inclusion of \mathcal{E}/U into \mathcal{E} .

The closed subtopos $\mathcal{E}_{\mathfrak{c}(U)}$ has reflector $K^* \colon \mathcal{E} \to \mathcal{E}_{\mathfrak{c}(U)}$ given on objects by the following pushout.



On morphisms $f: G \to G'$ in \mathcal{E} , $K^*(f)$ is given by the universal property of the pushout in the following diagram. Here the left, front and top faces commute and so a diagram chase determines that $p_1^{G'}f$ and $p_2^{G'} \mathrm{id}_U$ indeed form a cocone.



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We denote the right adjoint of K^* by K and write ζ and δ for the unit and counit respectively.

As expected, $\mathcal{E}_{\mathfrak{o}(U)}$ and $\mathcal{E}_{\mathfrak{c}(U)}$ are complemented subobjects. Given toposes \mathcal{H} and \mathcal{N} we can ask for which toposes \mathcal{G} may \mathcal{H} be embedded as an open subtopos and \mathcal{N} its closed complement. This is solved completely by the Artin glueing construction. For any finite-limit-preserving functor $F: \mathcal{H} \to \mathcal{N}$ we may construct the category $\operatorname{Gl}(F)$ whose objects are triples (N, H, ℓ) in which $N \in \mathcal{N}, H \in \mathcal{H}$ and $\ell: N \to F(H)$ and whose morphisms are pairs (f, g) making the following diagram commute.



The category $\operatorname{Gl}(F)$ is a topos, and moreover, the obvious projections $\pi_1: \operatorname{Gl}(F) \to \mathcal{N}$ and $\pi_2: \operatorname{Gl}(F) \to \mathcal{H}$ are finite-limit preserving. The projection π_2 has a right adjoint π_{2*} sending objects H to $(F(H), H, \operatorname{id}_{F(H)})$ and morphisms f to (F(f), f). This map π_{2*} is a geometric morphism and, in particular, an open subtopoly inclusion. Similarly, π_1 has a right adjoint sending objects N to (N, 1, !) and morphisms f to (f, !) where $!: N \to 1$ is the unique map to the terminal object. This is itself a geometric morphism, and indeed, a closed subtopoly inclusion.

Remarkably, Artin glueings may be viewed as both comma and cocomma objects in the category of toposes with finite-limit-preserving functors. We provide a proof of the latter in Section 4.

One sees immediately that $\pi_1\pi_{2*} = F$. This suggests a way to view any open or closed subtopos as corresponding to one in glueing form. If $K: \mathcal{G}_{\mathfrak{c}(U)} \to \mathcal{G}$ and $E_*: \mathcal{G}/U \to \mathcal{G}$ are respectively the inclusions of open and closed subtoposes, then there is a natural sense in which these maps correspond to $\pi_{1*}: \mathcal{G}_{\mathfrak{c}(U)} \to$ $\operatorname{Gl}(K^*E_*)$ and $\pi_{2*}: \mathcal{G}/U \to \operatorname{Gl}(K^*E_*)$ respectively. This fact is well known, though a new proof will be provided in Section 3.2.

We now note that the maps π_1 : $\operatorname{Gl}(F) \to \mathcal{N}$ and π_2 : $\operatorname{Gl}(F) \to \mathcal{H}$ are fibrations. By the above argument, these results apply equally to the inverse image maps of open and closed subtoposes. This is likely well known, though we were unable to find explicit mention of this in the literature.

Proposition 2.5. Let $F: \mathcal{H} \to \mathcal{N}$ be a finite-limit-preserving map between toposes. Then the projection $\pi_2: \operatorname{Gl}(F) \to \mathcal{H}$ is a fibration.

Proof. This follows from [18, Theorem 14], but it also not hard to show this directly. \Box

The result of [18, Theorem 14] also implies that the other projection π_1 is an *opfibration*, though we will now prove that it is also a fibration.

Proposition 2.6. Let $F: \mathcal{H} \to \mathcal{N}$ be a finite-limit-preserving map between toposes. Then the projection $\pi_1: \operatorname{Gl}(F) \to \mathcal{N}$ is a fibration.

Proof. We must show that if $f: N' \to N$ is a morphism in \mathcal{N} then for every object (N, H, ℓ) in Gl(F) there exists a cartesian lifting. This map is given by $(f, id_H): (N', H, \ell f) \to (N, H, \ell)$.

To see that this map satisfies the universal property, suppose that we have a morphism $(g_1, g_2): (A, B, k) \rightarrow (N, H, \ell)$ which is mapped by π_1 to fh. We must show there is a unique map \overline{h} such that $(g_1, g_2) = (f, \mathrm{id}_H)\overline{h}$ and $\pi_1(\overline{h}) = h$. These constraints imply that $\overline{h} = (h, g_2)$.

To see that (h, g_2) is a morphism in Gl(F), we consider the following diagram. The left-hand square commutes as $\pi_1(g_1, g_2) = fh$ and the right-hand square commutes since (g_1, g_2) is a morphism.



Finally, we immediately see that $(f, id_H)(h, g_2) = (g_1, g_2)$, as required. \Box

3. Adjoint extensions

In generalising the frame results to the topos setting, it is clear that the appropriate 2-category to consider is Top_{lex} , the 2-category of toposes, finite-limit-preserving functors and natural transformations. (This is the horizontal bicategory of the double category of toposes considered in [15].) For convenience, we will assume that 1 always refers to a distinguished terminal object in a topos, and 0 a distinguished initial object.

We will now introduce the necessary concepts in order to discuss extensions of toposes and show how Artin glueings can be viewed as adjoint extensions. In particular, the definition of extension will require notions of kernel and cokernel.

3.1. Zero morphisms, kernels and cokernels

The definition of extensions requires a notion of zero morphisms. Let us now define these in the 2categorical context.

Definition 3.1. A *pointed 2-category* is a 2-category equipped with a class \mathcal{Z} of 1-morphisms (called *zero morphisms*) satisfying the following conditions:

- \mathcal{Z} contains an object of each hom-category,
- \mathcal{Z} is an ideal with respect to composition (as in [9]) that is, $g \in \mathcal{Z} \implies fgh \in \mathcal{Z}$,
- \mathcal{Z} is closed under 2-isomorphism in the sense that if $f \in \mathcal{Z}$ and $f' \cong f$ then $f' \in \mathcal{Z}$,
- for any parallel pair f_1, f_2 of morphisms in \mathcal{Z} , there is a unique 2-morphism $\xi: f_1 \to f_2$.

Definition 3.2. A zero object in a 2-category is an object which is both 2-initial and 2-terminal.

Lemma 3.3. A 2-terminal or 2-initial object in a pointed 2-category is always a zero object. Furthermore, any 2-category with a zero object has a unique pointed 2-category structure where the zero morphisms are those 1-morphisms which factor through the zero object up to 2-isomorphism.

Remark 3.4. Definition 3.1 is a categorification of pointed-set-enriched categories. Pointed categories are often instead defined as those having a zero object. The previous lemma shows that our definition agrees with a definition in terms of zero objects when the category has a terminal or initial object.

Definition 3.5. The 2-cokernel of a morphism $f: A \to B$ in a pointed 2-category is an object C equipped with a morphism $c: B \to C$ such that cf is a zero morphism and which is the universal such in the following sense.

- (1) If $t: B \to X$ is such that tf is a zero morphism, then there exists a morphism $h: C \to X$ such that hc is isomorphic to t.
- (2) Given $h, h': C \to X$ and $\alpha: hc \to h'c$, there is a unique $\gamma: h \to h'$ such that $\gamma c = \alpha$.

The 2-kernel of a morphism in a pointed 2-category \mathcal{C} is simply the 2-cokernel in \mathcal{C}^{op} .

A similar notion was defined for groupoid-enriched categories in [8].

Note that these may also be defined in terms of 2-pushouts or 2-coequalisers involving the zero morphism (in the same way kernels and cokernels can be defined in the 1-categorical setting). Although there would usually be coherence conditions that need to be satisfied, the uniqueness of natural isomorphisms between zero morphisms eliminates all of them in the case of 2-cokernels.

Remark 3.6. Note that the condition (2) for 2-kernels is simply the statement that the 2-kernel map is a fully faithful 1-morphism. Moreover, since an adjoint of a fully faithful morphism is fully faithful in the opposite 2-category, we have that when a putative 2-cokernel $c: B \to C$ has a (left or right) adjoint d, then condition (2) for 2-cokernels is equivalent to d being fully faithful.

We can now consider how these concepts behave in our case of interest. Note that Top_{lex} has a zero object, the trivial topos. Then zero morphisms in Top_{lex} are precisely those functors which send every object to a terminal object.

In Top_{lex}, 2-cokernels of morphisms $F: \mathcal{N} \to \mathcal{G}$ always exist and are given by the open subtopos corresponding to F(0).

Proposition 3.7. The 2-cokernel of $F \colon \mathcal{N} \to \mathcal{G}$ is given by $E \colon \mathcal{G} \to \mathcal{G}/F(0)$ sending objects G to $G \times F(0)$ and morphisms $f \colon G \to G'$ to $(f, id_{F(0)}) \colon G \times F(0) \to G' \times F(0)$.

Proof. We know that E lies in Top_{lex} and so we begin by showing that EF is a zero morphism. The terminal object in $\mathcal{G}/F(0)$ is F(0) and so consider the following calculation.

$$EF(N) = F(N) \times F(0)$$
$$\cong F(N \times 0)$$
$$\cong F(0).$$

Next suppose that $T: \mathcal{G} \to \mathcal{X}$ is such that TF is a zero morphism. We claim that $TE_*: \mathcal{H} \to \mathcal{X}$ when composed with E is naturally isomorphic to T.

Observe that

$$T(G) \cong T(G) \times 1$$
$$\cong T(G) \times TF(0)$$
$$\cong T(G \times F(0))$$
$$\cong T(E_!E(G))$$

where each isomorphism is natural in G. Hence, $T \cong TE_!E \cong (TE_!E)E_*E \cong TE_*E$ where the central isomorphism comes from $E\theta \colon E \xrightarrow{\sim} EE_*E$.

The final condition for the 2-cokernel holds immediately because E has a full and faithful adjoint. \Box

Unfortunately, 2-kernels do not always exist in Top_{lex} . However, they do exist in the larger 2-category Cat_{lex} of finitely-complete categories and finite-limit-preserving functors.

Proposition 3.8. Let $F: \mathcal{G} \to \mathcal{H}$ be a morphism in $\operatorname{Cat}_{\operatorname{lex}}$. The 2-kernel of F, which we write as $\operatorname{Ker}(F)$, is given by the inclusion into \mathcal{G} of the full subcategory of objects sent by F to a terminal object.

Proof. Since F preserves finite limits and sends each object in Ker(F) to a terminal object, it is clear that Ker(F) is closed under finite limits. Naturally, the inclusion is a finite-limit-preserving functor.

It is clear that FK is a zero morphism. We must check that if $T: \mathcal{X} \to \mathcal{G}$ is such that FT is a zero morphism, then it factors through $\operatorname{Ker}(F)$. Note that since FT is a zero morphism, all objects (and morphisms) in its image lie in $\operatorname{Ker}(F)$. Thus, it is easy to see that T factors through $\operatorname{Ker}(F)$. The uniqueness condition of the universal property is immediate, as the inclusion of $\operatorname{Ker}(F)$ is full and faithful. \Box

We will only be concerned with 2-kernels of 2-cokernels. The following proposition shows that these do always exist in Top_{lex} .

Proposition 3.9. Let U be a subterminal of a topos \mathcal{G} and consider $E: \mathcal{G} \to \mathcal{G}/U$ defined as in Proposition 3.7. Then the kernel of E is given by $K: \mathcal{G}_{\mathfrak{c}(U)} \hookrightarrow \mathcal{G}$, the inclusion of the closed subtopos corresponding to U.

Proof. Since Top_{lex} is a full sub-2-category of Cat_{lex} , it suffices to show that the closed subtopos $\mathcal{G}_{\mathfrak{c}(U)}$ is equivalent to Ker(E), the full subcategory of objects sent by E to a terminal object.

The reflector $K^* \colon \mathcal{G} \to \mathcal{G}_{\mathfrak{c}(U)}$ sends an object G to the following pushout.



First we show that $K^*(G)$ lies in Ker(E). We know that E preserves colimits and so we obtain the following pushout in \mathcal{G}/U .



But note that p is an isomorphism, since it is the pushout of an identity morphism and thus $EK^*(G) \cong U$ and $K^*(G)$ lies in Ker(E).

Finally, we must show that K^* fixes the objects of $\operatorname{Ker}(E)$. First observe that U is the initial object in $\operatorname{Ker}(E)$, since if X is an object in $\operatorname{Ker}(E)$ then we have $\operatorname{Hom}_{\mathcal{G}}(U,X) = \operatorname{Hom}_{\mathcal{G}}(E_!(U),X) \cong \operatorname{Hom}_{\mathcal{G}/U}(U,E(X)) \cong \operatorname{Hom}_{\mathcal{G}/U}(U,U)$. There is precisely one morphism in $\operatorname{Hom}_{\mathcal{G}/U}(U,U)$, since U is the terminal object in \mathcal{G}/U .

Now consider the following candidate pushout diagram where G lies in Ker(E).



To see that the square commutes, note that by assumption G lies in Ker(E) and so $G \times U \cong U$. Therefore, $G \times U$ is initial in Ker(E) and there is a unique map into G.

Now suppose (f,g) is a cocone in \mathcal{G} . It is clear that the candidate morphism $h: G \to X$ must equal f and so we must just show that $f \circ !_G = g$. Since $G \times U$ and U are both initial in Ker(E), π_U has an inverse $!_{G \times U}: U \to G \times U$.

We now have

$$f \circ !_G = f!_G \pi_U !_{G \times U}$$
$$= f \pi_G !_{G \times U}$$
$$= g \pi_U !_{G \times U}$$
$$= g.$$

This gives that G is the pushout and hence fixed by K^* . \Box

3.2. Adjoint extensions and Artin glueings

We are now in a position to define our main object of study: adjoint split extensions.

Definition 3.10. A diagram in Top_{lex} of the form

$$\mathcal{N} \xrightarrow{K} \mathcal{G} \xleftarrow{E}_{K} \mathcal{H}$$

equipped with a natural isomorphism $\varepsilon \colon EE_* \to \operatorname{Id}_{\mathcal{H}}$ is called an *adjoint split extension* if K is the 2-kernel of E, E is the 2-cokernel of K, E_* is the right adjoint of E and ε is the counit of the adjunction.

Remark 3.11. Propositions 3.7 and 3.9 suggest that every adjoint split extension is equivalent to an extension arising from a closed subtopos and its open complement (in a sense that will be made precise in Definition 4.1). Here U = K(0) is a subterminal in \mathcal{G} .



The above situation is precisely the setting in which Artin glueings are studied and it is well known that in this case \mathcal{G} is equivalent to an Artin glueing $\operatorname{Gl}(K^*E_*)$. We will present an alternative proof of this result from the perspective of extensions.

We begin by showing that Artin glueings can be viewed as adjoint split extensions in a natural way.

Proposition 3.12. Let $F: \mathcal{H} \to \mathcal{N}$ be a finite-limit-preserving functor. Then the diagram

$$\mathcal{N} \xrightarrow{\pi_{1*}} \operatorname{Gl}(F) \xrightarrow{\pi_2} \mathcal{H}$$

is an adjoint split extension in Top_{lex}.

Proof. We first note that Gl(F) is a topos. This is a fundamental result in the theory of Artin glueings of toposes and a proof can be found in [19].

By Proposition 3.8, it is immediate that π_{1*} is the 2-kernel of π_2 . To see that π_2 is the 2-cokernel of π_{1*} , we first observe that the slice category of $\operatorname{Gl}(F)$ by the subterminal $(0, 1, !) = \pi_{1*}(0)$ is equivalent to \mathcal{H} . The objects of $\operatorname{Gl}(F)/(0, 1, !)$ are isomorphic to those the form (0, H, !) (since every morphism into an initial object in \mathcal{N} is an isomorphism) and its morphisms of the form (!, f). If $L: \operatorname{Gl}(F)/(0, 1, !) \to \mathcal{H}$ is this isomorphism sending (0, H, !) to H and (!, f) to f and $E: \operatorname{Gl}(F) \to \operatorname{Gl}(F)/(0, 1, !)$ is the cokernel map, then it is clear that $LE \cong \pi_2$. \Box

The following proposition is shown for Grothendieck toposes in [1], but deserves to be more well known. Here we prove it for general elementary toposes. (It also follows easily from the theory of Artin glueings, but here we will use it to develop that theory.)

Recall that we use θ for the unit of the open subtopos adjunction $E \dashv E_*$ and ζ for the unit of the closed subtopos adjunction $K^* \dashv K$.

Proposition 3.13. Let \mathcal{G} be a topos and consider an open subtopos $E_* \colon \mathcal{H} \hookrightarrow \mathcal{G}$ with closed complement $K \colon \mathcal{N} \hookrightarrow \mathcal{G}$. Then each object G in \mathcal{G} can be expressed as the following pullback in \mathcal{G} of objects from \mathcal{N} and \mathcal{H} .



Proof. First note that the diagram commutes by the naturality of ζ . Recall that, setting $U = KK^*(0) \cong K(0)$, we have $E_*E(G) = (G \times U)^U \cong G^U$ and that $KK^*(G)$ is the pushout of G and U along the projections $\pi_1 \colon G \times U \to G$ and $\pi_2 \colon G \times U \to U$. Now we can rewrite the relevant pullback diagram as follows.



Here the ι maps are injections into the pushout and c is the unit of the exponential adjunction, which intuitively maps elements of G to their associated constant functions.

Let us express P in the internal logic. We have

$$P = \{(f,[g]) \mid f \sim c(g)\} \cup \{(f,[*]) \mid f \sim *\} \subseteq G^U \times (G+U)/\sim$$

where \sim denotes the equivalence relation generated by $f \sim *$ for $* \in U$. Explicitly, we find that $f \sim f' \iff f = f' \lor * \in U$. Thus, we find

$$P = \{ (f, [g]) \mid f = c(g) \lor * \in U \} \cup \{ (f, [*]) \mid f \in G^U, * \in U \}$$

= $\{ (c(g), [g]) \mid g \in G \} \cup \{ (f, [g]) \mid * \in U \} \cup \{ (f, [*]) \mid * \in U \}.$

Now observe that if $* \in U$, then [g] = [*] and hence $\{(f, [g]) | * \in U\} \subseteq \{(f, [*]) | * \in U\}$. Finally, commuting the subobject and the quotient we arrive at

$$P = \left(\{ (c(g), g) \mid g \in G \} \sqcup \{ (f, *) \mid * \in U \} \right) / \sim$$

where the equivalence relation is generated by $(f,g) \sim (f,*)$ for $* \in U$. Note that the union is now disjoint.

The map $r: G \to P$ sends $g \in G$ to [(c(g), g)]. We can define a candidate inverse by $s: P \to G$ by $(f,g) \mapsto g$ and $(f,*) \mapsto f(*)$, which can seen to be well-defined, since if (f,g) and (f',*) are elements of the disjoint union with f = f' and $* \in U$, then f'(*) = c(g)(*) = g.

We clearly have $sr = id_G$. Now if $[(c(g), g)] \in P$ then rs([(c(g), g)]) = r(g) = [(c(g), g)]. On the other hand, if $[(f, *)] \in P$ then $* \in U$ and rs([f, *]) = r(f(*)) = [(c(f(*)), *)], which equals [(f, *)] since f(*) = c(f(*))(*) so that f = c(f(*)). Thus, r and s are inverses as required. \Box

Remark 3.14. The non-classical logic in the above proof can be hard to make sense of. It can help to consider the cases where U = 0 and U = 1. In the former case, G^U contains no information and $G +_{G \times U} U \cong G$, while in the latter case the opposite is true. The general case 'interpolates' between these.

The next result will play a central role in this paper.

Proposition 3.15. Let \mathcal{G} be a topos and consider an open subtopos $E_* \colon \mathcal{H} \hookrightarrow \mathcal{G}$ with closed complement $K \colon \mathcal{N} \hookrightarrow \mathcal{G}$. We have the following pullback in the category of (finite-limit-preserving) endofunctors on \mathcal{G} .



Proof. We proved this on objects in Proposition 3.13. Now consider a morphism $p: G \to G'$ in \mathcal{G} . We obtain the following commutative cube.



An enjoyable diagram chase around the cube shows that p is the unique morphism making the cube commute by the pullback property of the front face. Since pullbacks in the functor category are computed pointwise, this yields the desired result. \Box

Remark 3.16. It is remarked in [11] that adjoint extensions of frames can be viewed as weakly Schreier split extensions of monoids as defined in [4]. Propositions 3.13 and 3.15 can be viewed as a categorified version of the weakly Schreier condition for the topos setting, though it is as yet unclear how the weakly Schreier condition might be categorified more generally. It would also be interesting to see how the general theory of (S-)protomodular categories [2,6] might be categorified. Another potential example of 2-dimensional protomodularity can be found in [12].

We can now prove the main result of this section. For now we will treat equivalences of adjoint extensions without worrying too much about coherence, which we will discuss in more detail when we define morphisms of extensions in Section 4.

Theorem 3.17. Let $\mathcal{N} \xrightarrow{K} \mathcal{G} \xrightarrow{E} \mathcal{H}$ be an adjoint split extension. Then it is equivalent to the split extension $\mathcal{N} \xrightarrow{\pi_{1*}} \operatorname{Gl}(K^*E_*) \xrightarrow{\pi_2} \mathcal{H}.$

Proof. We denote the unit of $\pi_2 \dashv \pi_{2*}$ by θ' : $\mathrm{Id}_{\mathrm{Gl}(K^*E_*)} \to \pi_{2*}\pi_2$ and observe that its component at (N, H, ℓ) can be given explicitly by (ℓ, id_H) : $(N, H, \ell) \to (K^*E_*(H), H, \mathrm{id}_H)$. We denote the unit of $\pi_1 \dashv \pi_{1*}$ by ζ' : $\mathrm{Id}_{\mathrm{Gl}(K^*E_*)} \to \pi_{1*}\pi_1$ and its components are given by $(\mathrm{id}_N, !)$: $(N, H, \ell) \to (N, 1, !)$.

Consider the functor $\Phi: \mathcal{G} \to \operatorname{Gl}(K^*E_*)$ which sends objects G to the triple $(K^*(G), E(G), K^*\theta_G)$ and morphisms $f: G \to G'$ to $(K^*(f), E(f))$. This is a morphism of extensions in the sense that we have natural isomorphisms $\alpha: \pi_{1*} \to \Phi K$, $\beta: \pi_2 \Phi \to E$ and $\gamma: \pi_{2*} \to \Phi E_*$ given by $\alpha = (\operatorname{id}, !) \circ \pi_{1*} \delta^{-1}$ (where $\delta: K^*K \to \operatorname{Id}_{\mathcal{N}}$ is the counit of $K^* \dashv K$ and $(\operatorname{id}, !): \pi_{1*}K^*K \cong \Phi K$ is simply to ensure the different choices of terminal object agree), $\beta = \operatorname{id}$ (using $\pi_2 \Phi = E$) and $\gamma = (\operatorname{id}_{K^*E_*}, \varepsilon^{-1})$ (where $\varepsilon: EE_* \to \operatorname{Id}_{\mathcal{H}}$ is the counit of $E \dashv E_*$ and this is a morphism in $\operatorname{Gl}(K^*E_*)$ by the triangle identity). This can be seen to be a morphism of extensions as defined in Definition 4.1 (see Remark 4.2).

We must show that Φ is an equivalence. We claim that the following pullback in Hom $(Gl(K^*E_*), \mathcal{G})$ is the inverse of Φ . (Here the bottom equality comes from $K^*E_* = \pi_1\pi_{2*}$.)



We shall make extensive use of Proposition 3.15 in order to prove this.

To see that $\Phi'\Phi \cong id_{\mathcal{G}}$ note that composition with Φ on the right preserves limits and thus can be represented as the following pullback in $\operatorname{Hom}(\mathcal{G}, \mathcal{G})$.



Note that $\pi_1 \Phi = K^*$, $\pi_2 \Phi = E$ and that $\theta'_{\Phi(G)} = (K^* \theta_G, \operatorname{id}_{E(G)})$ which of course gives that $K \pi_1 \theta' \Phi = KK^* \theta$. After making these substitutions into the diagram above, we have the pullback square occurring in Proposition 3.15, which by the universal property gives that $\Phi \Phi'$ is naturally isomorphic to $\operatorname{Id}_{\mathcal{G}}$.

The same idea works for $\Phi\Phi'$. We consider the following pullback square in the category of endofunctors on $\operatorname{Gl}(K^*E_*)$.



The bottom trapezium commutes by naturality of α^{-1} (and using $K^*E_* = \pi_1\pi_{2*}$), whereas both paths around the right-hand trapezium can be seen to compose to $(\mathrm{id}_{K^*E_*(H)}, !)$. Thus this pullback diagram has the same form as that in Proposition 3.15 and so Proposition 3.12 allows us to deduce that $\Phi\Phi'$ is naturally isomorphic to the identity, completing the proof. \Box

Together with Proposition 3.12 this shows that adjoint split extensions and Artin glueings are essentially the same. The equivalence Φ so defined is natural in a sense that will become clear later in Section 4.2.

4. The category of extensions

It follows from Theorem 3.17 that equivalence classes of adjoint extensions between \mathcal{H} and \mathcal{N} are in bijection with those of Hom $(\mathcal{H}, \mathcal{N})$. However, the extensions have a natural 2-categorical structure and so we would like to know how this relates to the categorical structure of Hom $(\mathcal{H}, \mathcal{N})$.

4.1. Morphisms of extensions

Definition 4.1. Suppose we have two adjoint extensions,

$$\mathcal{N} \xrightarrow{K_1} \mathcal{G}_1 \xleftarrow{E_1}{\longleftarrow} \mathcal{H} \text{ and } \mathcal{N} \xrightarrow{K_2} \mathcal{G}_2 \xleftarrow{E_2}{\longleftarrow} \mathcal{H},$$

with the same kernel and cokernel objects and associated isomorphisms ε_1 and ε_2 respectively. Consider the following diagram where α , β and γ are natural isomorphisms and Ψ is a finite-limit-preserving functor.



We say the functor $\Psi: \mathcal{G}_1 \to \mathcal{G}_2$ together with α , β and γ is a morphism of adjoint extensions (of \mathcal{H} by \mathcal{N}) if $\varepsilon_2 = \varepsilon_1(\beta E_{1*})(E_2\gamma)$.

Given two such morphisms $(\Psi, \alpha, \beta, \gamma)$ and $(\Psi', \alpha', \beta', \gamma')$, a 2-morphism of adjoint extensions is a natural transformation $\tau: \Psi \to \Psi'$ such that $\alpha' = (\tau K_1)\alpha$, $\beta = \beta'(E_2\tau)$ and $\gamma' = (\tau E_{1*})\gamma$.

The morphisms compose in the obvious way: by composing the functors and pasting the natural transformations together by juxtaposing the squares from the diagram above. Horizontal and vertical composition of 2-morphisms is given by the corresponding operations of the natural transformations τ . It is not hard to see that this gives a strict 2-category $\text{Ext}(\mathcal{H}, \mathcal{N})$.

Remark 4.2. Note that the 'equivalence of extensions' $(\Phi, \alpha, \beta, \gamma)$ defined in Theorem 3.17 is indeed a morphism of extensions in the above sense, since there ε_2 and β are identities and $\varepsilon(\pi_2\gamma) = \varepsilon\varepsilon^{-1} = id$, as required.

Lemma 4.3. Given two adjoint extensions as above, a functor $\Psi: \mathcal{G}_1 \to \mathcal{G}_2$ and natural transformations $\alpha: K_2 \to \Psi K_1$ and $\gamma: E_{2*} \to \Psi E_{1*}$, there is a unique natural isomorphism $\beta: E_2 \Psi \to E_1$ making $(\Psi, \alpha, \beta, \gamma)$ a morphism of adjoint extensions.

Furthermore, given two morphisms of adjoint extensions as above, a natural transformation $\tau: \Psi \to \Psi'$ is a 2-morphism of adjoint extensions if and only if $\alpha' = (\tau K_1)\alpha$ and $\gamma' = (\tau E_{1*})\gamma$.

Proof. Any such β must satisfy $\varepsilon_2 = \varepsilon_1(\beta E_{1*})(E_2\gamma)$. But this can be rewritten as $\varepsilon_2(E_2\gamma^{-1}) = \varepsilon_1(\beta E_{1*})$, which shows that β and γ^{-1} are mates with respect to the adjunctions $E_1 \dashv E_{1*}$ and $E_2 \dashv E_{2*}$ and hence determine each other.

We now show that β so defined is an isomorphism. As the mate of γ^{-1} , we can express β as $(\varepsilon_2 E_1)(E_2\gamma^{-1}E_1)(E_2\Psi\theta_1)$. Now since ε_2 and γ^{-1} are isomorphisms, we need only show $E_2\Psi\theta_1$ is an isomorphism. This map occurs in the pullback obtained by applying $E_2\Psi$ to the pullback square of Proposition 3.15.



Now observe that $E_2\Psi K_1 \cong E_2K_2 \cong 1$ is a zero morphism and hence so are $E_2\Psi K_1K_1^*$ and $E_2\Psi K_1K_1^*E_{1*}E_1$. Therefore, the bottom arrow of the above diagram is an isomorphism, and as the pullback of an isomorphism, $E_2\Psi\theta_1$ is an isomorphism too.

Finally, we show that the condition on β for 2-morphisms of extensions is automatic. Simply observe the following string diagrams.



Here the first diagram represents $\beta'(E_2\tau)$ and the last diagram represents β . In moving from the first diagram to the second we shift τ above θ_1 and to move from the second diagram to the third we use $\gamma' = (\tau E_{1*})\gamma$. \Box

Lemma 4.4. Suppose $(\Psi, \alpha, \beta, \gamma)$ and $(\Psi', \alpha', \beta', \gamma')$ are parallel morphisms of extensions. Then any 2-morphism τ between them is unique and invertible.

Moreover, such a 2-morphism exists if and only if $\alpha' \alpha^{-1} K_1^* E_{1*} \circ \Psi \zeta_1 E_{1*} = \Psi' \zeta_1 E_{1*} \circ \gamma' \gamma^{-1}$.

Proof. Suppose $\tau: \Psi \to \Psi'$ is a 2-morphism of extensions. Then we have $\tau K_1 = \alpha' \alpha^{-1}$ and $\tau E_{1*} = \gamma' \gamma^{-1}$. Now by composing the pullback square of Proposition 3.15 with Ψ and Ψ' and using the naturality of τ we have the following commutative cube in $\operatorname{Hom}_{\operatorname{Top}_{\operatorname{lex}}}(\mathcal{G}_1, \mathcal{G}_2)$.



The universal property of the pullback on the front face then gives that τ is uniquely determined by $\tau K_1 K_1^*$ and $\tau E_{1*}E_1$, and hence by $\tau K_1 = \alpha' \alpha^{-1}$ and $\tau E_{1*} = \gamma' \gamma^{-1}$. Thus, the morphism τ is unique if it exists.

We can also attempt to use a similar cube to construct τ without assuming it exists a priori by replacing τK_1 with $\alpha' \alpha^{-1}$ and $\tau E_{1*} = \gamma' \gamma^{-1}$ in the above diagram. However, in order to obtain a map τ from the universal property of the pullback, we require that the right-hand face commutes.



Note that this square commutes if and only if the similar square obtained by inverting the horizontal morphisms commutes. But this latter square is precisely the square we need to commute to obtain a 2-morphism in the opposite direction. Uniqueness then shows that these two 2-morphisms compose to give identities.

Finally, observe that commutativity of this square is the required equality stated above whiskered with E_1 . This is equivalent to the desired condition, since E_1 is essentially surjective. \Box

Corollary 4.5. The 2-category of adjoint extensions $\mathbb{E}xt(\mathcal{H}, \mathcal{N})$ is equivalent to the locally trivial 2-category $\mathbb{E}xt(\mathcal{H}, \mathcal{N})$ of adjoint extensions and isomorphism classes of morphisms.

This will justify working with $Ext(\mathcal{H}, \mathcal{N})$ (viewed as a 1-category) going forward.

Lemma 4.6. From a morphism of extensions $(\Psi, \alpha, \beta, \gamma)$ we can form an associated natural transformation from $K_2^*E_{2*}$ to $K_1^*E_{1*}$ given by $(\delta_2 K_1^*E_{1*})(K_2^*\alpha^{-1}K_1^*E_{1*})(K_2^*\Psi\zeta_1E_{1*})(K_2^*\gamma)$ where δ_2 is the counit of the $K_2^* \dashv K_2$ adjunction and which is depicted below. Two parallel morphisms of extensions are isomorphic if and only if their corresponding natural transformations are equal.



Proof. There is a (necessarily invertible) 2-morphism from $(\Psi, \alpha, \beta, \gamma)$ to $(\Psi', \alpha', \beta', \gamma')$ if and only if $(\alpha' \alpha^{-1} K_1^* E_{1*})(\Psi \zeta_1 E_{1*}) = (\Psi' \zeta_1 E_{1*})\gamma' \gamma^{-1}$. We can now move all the unprimed variables to the left and primed variables to the right by multiplying both sides of this equation on the left by $\alpha'^{-1} K_1^* E_{1*}$ and on the right by γ to obtain $(\alpha^{-1} K_1^* E_{1*})(\Psi \zeta_1 E_{1*})\gamma = (\alpha'^{-1} K_1^* E_{1*})(\Psi' \zeta_1 E_{1*})\gamma'$. These are the mates of the desired natural transformations with respect to the adjunction $K_2^* \dashv K_2$. \Box

4.2. The equivalence of categories

In this section we show that the categories $\text{Ext}(\mathcal{H}, \mathcal{N})$ and $\text{Hom}(\mathcal{H}, \mathcal{N})^{\text{op}}$ are equivalent. This requires showing that isomorphism classes of morphisms of extensions correspond to natural transformations. We have already seen that each isomorphism class has an associated natural transformation. We will now further explore this relationship, making use of the following folklore result.

Proposition 4.7. Let $\mathcal{N} \xrightarrow{K} \mathcal{G} \xleftarrow{E}_{E_*} \mathcal{H}$ be an adjoint extension. Then the following diagram is a cocomma square.



Proof. We first check the 2-categorical condition. Consider two finite-limit-preserving functors $U, V: \mathcal{G} \to \mathcal{X}$ and natural transformations $\mu: UE_* \to VE_*$ and $\nu: UK \to VK$ such that $(V\zeta E_*)\mu = (\nu K^*E_*)(U\zeta E_*)$. We must find a unique $\omega: U \to V$ such that $\omega E_* = \mu$ and $\omega K = \nu$.

We use Proposition 3.15 to express U and V as pullbacks and then as in Lemma 4.4 we find that there is a unique map $\omega: U \to V$ with $\omega E_* = \mu$ and $\omega K = \nu$ as long as the following diagram commutes.



But commutativity of this diagram is simply the assumed condition whiskered with E on the right.

Now we show the 1-categorical condition. Suppose we have finite-limit-preserving functors $T_1: \mathcal{H} \to \mathcal{X}$ and $T_2: \mathcal{N} \to \mathcal{X}$ and a natural transformation $\varphi: T_1 \to T_2 K^* E_*$. We must construct a finite-limit-preserving functor $L: \mathcal{G} \to \mathcal{X}$ and natural isomorphisms $\tau_1: LE_* \to T_1$ and $\tau_2: LK \to T_2$ depicted below such that $\varphi = \tau_2 K^* E_* \circ L \zeta E_* \circ \tau_1^{-1}$.



Suppose we are given such a functor L and natural isomorphisms τ_1, τ_2 and consider the following pullback diagram.



Here the bottom trapezium commutes by the naturality of $\tau_2 K^*$ and the right trapezium commutes since $(\varphi \tau_1)E = (\tau_2 K^* E_* \circ L \zeta E_*)E$ by assumption. Note that the left edge of the large square is the mate of τ_2 with respect to $K^* \dashv K$ and the top edge is the mate of τ_1 with respect to $E \dashv E_*$.

Now without assuming L exists to start with, we can use the outer pullback diagram to define it and we may recover τ_1 and τ_2 as the mates of the resulting pullback projections.

Observe that precomposing the pullback with K turns the right-hand edge into an isomorphism between zero morphisms. Hence the left-hand morphism $(\tau_2 K^*)(L\zeta)K$ is an isomorphism as well. Since τ_2 is given by composing this with the isomorphism $T_2\delta$, we find that τ_2 is an isomorphism. On the other hand, precomposing the pullback with E_* turns the bottom edge into an isomorphism (as $\mathcal{N} \stackrel{E_*}{\longrightarrow} \mathcal{G}$ is a reflective subcategory). It follows that τ_1 is also an isomorphism.

Finally, we show that φ be can recovered in the appropriate way. The commutativity of the pullback square gives $\varphi E \circ \tau_1 E \circ L\theta = T_2 K^* \theta \circ \tau_2 K^* \circ L\zeta$. The result of whiskering this with E_* on the right and composing with $T_2 K^* E_* \varepsilon$ is depicted in the string diagram below.



The desired equality follows after using the triangle identities to 'pull the wires straight'. \Box

Proposition 4.8. Let A and B denote the adjoint extensions $\mathcal{N} \xrightarrow{K_1} \mathcal{G}_1 \xrightarrow{E_1} \mathcal{H}$ and

 $\mathcal{N} \xrightarrow{K_2} \mathcal{G}_2 \xrightarrow{E_2} \mathcal{H} \text{ respectively. There is a bijection between the equivalence classes of morphisms of extensions Hom(A, B) and the natural transformations Hom(K_2^*E_{2*}, K_1^*E_{1*}).$

Proof. Let $\psi: K_2^* E_{2*} \to K_1^* E_{1*}$ be a natural transformation and consider its mate natural transformation $\overline{\psi}: E_{2*} \to K_2 K_1^* E_{1*}$ with respect to the adjunction $K_2^* \dashv K_2$. Notice that $\overline{\psi}$ determines a 2-cocone in the following cocomma object diagram.



By the universal property of the cocomma, we get a map $\Psi: \mathcal{G}_1 \to \mathcal{G}_2$ and natural isomorphisms $\alpha: K_2 \to \Psi K_1$ and $\gamma: E_{2*} \to \Psi E_{1*}$. By Lemma 4.3 we can derive a unique natural isomorphism $\beta: E_2 \Psi \to E_1$ such that $(\Psi, \alpha, \beta, \gamma)$ is a morphism of extensions.

For the other direction we begin with a morphism of extensions $(\Psi, \alpha, \beta, \gamma)$ and form the pasting diagram above. We may consider the composite natural transformation

$$\overline{\psi} = \alpha^{-1} K_1^* E_{1*} \circ \Psi \zeta E_{1*} \circ \gamma \colon E_{2*} \to K_2 K_1^* E_{1*}.$$

Again we may use the adjunction $K_2^* \dashv K_2$ to arrive at the natural transformation

$$\delta_2 K_1^* E_{1*} \circ K_2^* \alpha^{-1} K_1^* E_{1*} \circ K_2^* \Psi \zeta E_{1*} \circ K_2^* \gamma \colon K_2^* E_{2*} \to K_1^* E_{1*}$$

where $\delta_2 \colon K_2^* K_2 \to \mathrm{Id}_{\mathcal{N}}$ is the counit of the adjunction.

It is clear that these processes are inverses by the uniqueness of the universal property (bearing in mind that the morphisms in $\text{Ext}(\mathcal{H}, \mathcal{N})$ are isomorphism classes). \Box

Remark 4.9. Notice that the natural transformation associated to a morphism of extensions $(\Psi, \alpha, \beta, \gamma)$ in the above proof is precisely the one described in Lemma 4.6.

Corollary 4.10. Let $F_1, F_2: \mathcal{H} \to \mathcal{N}$ be finite-limit-preserving functors and let

$$\Gamma_1 = \mathcal{N} \xrightarrow{\overset{} \longrightarrow } \operatorname{Gl}(F_1) \xrightarrow{\overset{} \pi_2^{F_1}} \mathcal{H} \quad and \quad \Gamma_2 = \mathcal{N} \xrightarrow{\overset{} \longrightarrow } \operatorname{Gl}(F_2) \xrightarrow{\overset{} \pi_2^{F_2}} \mathcal{H}$$

be the corresponding glueing extensions. Then $\operatorname{Hom}(F_2, F_1) \cong \operatorname{Hom}(\Gamma_1, \Gamma_2)$.

Proof. Since $\pi_1^{F_1}\pi_{2*}^{F_1} = F_1$ and $\pi_1^{F_2}\pi_{2*}^{F_2} = F_2$, the above proposition implies that $\operatorname{Hom}(\Gamma_1, \Gamma_2) \cong \operatorname{Hom}(F_2, F_1)$, as required. \Box

We are now ready to show that the categories $\text{Ext}(\mathcal{H}, \mathcal{N})$ and $\text{Hom}(\mathcal{H}, \mathcal{N})^{\text{op}}$ are equivalent. (A similar result can also be obtained from the results of [17].)

Definition 4.11. Let $\Gamma_{\mathcal{H},\mathcal{N}}$: Hom $(\mathcal{H},\mathcal{N})^{\mathrm{op}} \to \mathrm{Ext}(\mathcal{H},\mathcal{N})$ be the functor sending $F: \mathcal{H} \to \mathcal{N}$ to the extension $\mathcal{N} \xrightarrow{\pi_{1*}} \mathrm{Gl}(F) \xleftarrow{\pi_{2*}} \mathcal{H}$ and sending natural transformations to the associated morphism of extensions described in Corollary 4.10.

Theorem 4.12. The functor $\Gamma_{\mathcal{H},\mathcal{N}}$: Hom $(\mathcal{H},\mathcal{N})^{\mathrm{op}} \to \mathrm{Ext}(\mathcal{H},\mathcal{N})$ is a part of an equivalence.

An inverse $\Gamma_{\mathcal{H},\mathcal{N}}^{-1}$ sends an extension $\mathcal{N} \xrightarrow{K} \mathcal{G} \xleftarrow{E}_{E_*} \mathcal{H}$ to K^*E_* and a morphism of extensions to the natural transformation described in Lemma 4.6 and Proposition 4.8. For the adjunction $\Gamma_{\mathcal{H},\mathcal{N}}^{-1} \dashv \Gamma_{\mathcal{H},\mathcal{N}}$ we take the counit to be the identity and the unit to be given by the isomorphisms described in Theorem 3.17.

Proof. Note that $\Gamma_{\mathcal{H},\mathcal{N}}^{-1}\Gamma_{\mathcal{H},\mathcal{N}} = \mathrm{Id}_{\mathrm{Hom}(\mathcal{H},\mathcal{N})^{\mathrm{op}}}$. We see that $\Gamma_{\mathcal{H},\mathcal{N}}^{-1} \dashv \Gamma_{\mathcal{H},\mathcal{N}}$ with the identity as the counit, since for each natural transformation $\psi \colon F \to \Gamma^{-1}(A)$, there is a unique map $\Psi \colon A \to \Gamma_{\mathcal{H},\mathcal{N}}(F)$ such that $\Gamma_{\mathcal{H},\mathcal{N}}^{-1}(\Psi) = \psi$, namely the image of ψ under the inverse of the bijection

$$\operatorname{Hom}(A, \Gamma_{\mathcal{H}, \mathcal{N}}(F)) \cong \operatorname{Hom}(\Gamma_{\mathcal{H}, \mathcal{N}}^{-1}\Gamma_{\mathcal{H}, \mathcal{N}}(F), \Gamma_{\mathcal{H}, \mathcal{N}}^{-1}(A)) = \operatorname{Hom}(F, \Gamma_{\mathcal{H}, \mathcal{N}}^{-1}(A))$$

from Proposition 4.8.

Let A denote the extension $\mathcal{N} \xrightarrow{K} \mathcal{G} \xleftarrow{E}{E_*} \mathcal{H}$. It remains to show that the isomorphism of extensions $\Phi = (\Phi, \alpha, \beta, \gamma) \colon A \to \Gamma_{\mathcal{H}, \mathcal{N}}(K^*E_*)$ described in Theorem 3.17 is the component of the unit at A.

It suffices to show that the isomorphism $\operatorname{Hom}(A, \Gamma_{\mathcal{H},\mathcal{N}}\Gamma_{\mathcal{H},\mathcal{N}}^{-1}(A)) \cong \operatorname{Hom}(\Gamma_{\mathcal{H},\mathcal{N}}^{-1}(A), \Gamma_{\mathcal{H},\mathcal{N}}^{-1}(A))$ maps Φ to the identity — that is, that $\Gamma_{\mathcal{H},\mathcal{N}}^{-1}(\Phi) = \operatorname{id}_{K^*E_*}$. As in Lemma 4.6 we have $\Gamma_{\mathcal{H},\mathcal{N}}^{-1}(\Phi) = \delta' K^* E_* \circ \pi_1 \alpha^{-1} K^* E_* \circ \pi_1 \Phi \zeta E_* \circ \pi_1 \gamma$. Now recall that $\delta' : \pi_1 \pi_{1*} \to \operatorname{Id}_{\mathcal{N}}$ is the identity, $\alpha = (\operatorname{id}, !) \circ \pi_{1*} \delta^{-1}$ and $\gamma = (\operatorname{id}_{K^*E_*}, \varepsilon^{-1})$. Thus the expression can be seen to simplify to $\delta K^* E_* \circ K^* \zeta E_*$, which in turn is $\operatorname{id}_{K^*E_*}$ by the triangle identity for $K^* \dashv K$.

Finally, $\Gamma_{\mathcal{H},\mathcal{N}}$ and $\Gamma_{\mathcal{H},\mathcal{N}}^{-1}$ form an equivalence since the unit and counit are isomorphisms. \Box

With this in mind we may now consider the full subcategory of $\operatorname{Ext}(\mathcal{H}, \mathcal{N})$ whose objects are only those extensions of the form $\mathcal{N} \xrightarrow{\pi_{1*}} \operatorname{Gl}(F) \xleftarrow{\pi_{2*}} \mathcal{H}$ for some $F \colon \mathcal{H} \to \mathcal{N}$. It is evident that this full subcategory is equivalent to $\operatorname{Ext}(\mathcal{H}, \mathcal{N})$ and for the remainder of the paper we will choose to perform calculations in this subcategory for simplicity. We will discuss how this can be done coherently when we investigate the Ext 2-functor in Section 5.

We can now give a concrete description of the behaviour of morphisms of extensions. Suppose that $(\Psi, \alpha, \beta, \gamma) \colon \Gamma_{\mathcal{H}, \mathcal{N}}(F_1) \to \Gamma_{\mathcal{H}, \mathcal{N}}(F_2)$ is a morphism of adjoint extensions as in the following diagram.



Since $\alpha_N \colon (N, 1, !) \to \Psi(N, 1, !)$ is an isomorphism, we have $\Psi(N, 1, !) = (\overline{N}, \overline{1}, !)$ for some $\overline{N} \cong N$ and some terminal object $\overline{1}$. For simplicity, we will assume $\overline{1} = 1$ without any loss of generality. Similarly $\gamma_H \colon (F_2(H), H, \mathrm{id}) \to \Psi(F_1(H), H, \mathrm{id})$ is an isomorphism. So if $\Psi(F_1(H), H, \mathrm{id}) = (N_H, \overline{H}, t_H)$ then $t_H \colon N_H \to F_2(\overline{H})$ is an isomorphism.

Since Ψ preserves finite limits, we can use Proposition 3.13 to completely determine its behaviour. Every object in $\operatorname{Gl}(F_1)$ can be written as the following pullback diagram, where objects in the category are represented by the green arrows pointing out of the page and the pullback symbol has elongated into a wedge.



Note that the front and back faces are pullback squares in \mathcal{N} and the other faces correspond to morphisms in $\operatorname{Gl}(F_1)$.

We may now study how Ψ acts on this pullback diagram. Observe that the bottom face corresponds to the morphism $\pi_{1*}^{F_1}(\ell)$ where $\ell: N \to F_1(H)$. It is then sent by Ψ to the morphism represented in the diagram below.



The right face is the unit $\zeta_{\pi_{2*}^{F_1}(H)} \colon \pi_{2*}^{F_1}(H) \to \pi_{1*}^{F_1}F_1(H)$. Thus we have that Ψ sends this face to the following commutative square.



The pullback of these two faces will then give the image of (N, H, ℓ) under Ψ . The pullback diagram is given by the large cuboid in the diagram below. Here we have factored this pullback as in the similar pullback diagram in the proof of Proposition 4.7.



The bottom face of the bottom left cube and the right face of the top right cube are the commutative squares considered above. These have also been extended by the identity maps in the bottom right cube, so that the bottom and right-hand faces of the full cuboid are as required for the pullback in question.

The bottom left cube commutes by the naturality of α , while the top right cube commutes by the definition of $\psi = \Gamma_{\mathcal{H},\mathcal{N}}^{-1}(\Psi)$ as in Proposition 4.8. Since α and γ are isomorphisms, the top left cube is also a

pullback. Recall that the front and back faces are then also pullbacks. Since the top face of the top left cube must commute, we find that the green arrow we seek is given by $p_{\psi_H(\ell)}$. Hence, $\Psi(N, H, \ell)$ is isomorphic to $(N \times_{F_1(H)} F_2(H), H, p_{\psi_H}(\ell))$.

Of course, every natural transformation $\psi: F_2 \to F_1$ yields a morphism of extensions defined by such a pullback. For the associated natural isomorphisms we may take β to be the identity and α to be $(\hat{\alpha}, id)$ where $\hat{\alpha}$ is defined by the diagram below.



Finally, we take $\gamma = (\hat{\gamma}, id)$ where $\hat{\gamma}$ is specified by the diagram below.



It is easy to see that $\varepsilon_2 = \varepsilon_1(\beta \pi_{2*}^{F_1})(\pi_2^{F_2}\gamma)$ as each factor is just the identity.

We now end this section with what is perhaps a surprising result about morphisms of extensions.

Proposition 4.13. If $(\Psi, \alpha, \beta, \gamma)$: $\Gamma_{\mathcal{H}, \mathcal{N}}(F_1) \rightarrow \Gamma_{\mathcal{H}, \mathcal{N}}(F_2)$ is a morphism of adjoint extensions, then Ψ : $\operatorname{Gl}(F_1) \rightarrow \operatorname{Gl}(F_2)$ is a geometric morphism of toposes.

Proof. Let $\psi: F_2 \to F_1$ be the natural transformation associated to Ψ . We can construct a functor $\Psi^*: \operatorname{Gl}(F_2) \to \operatorname{Gl}(F_1)$ which sends (N, H, ℓ) to $(N, H, \psi_H \ell)$ and leaves morphisms 'fixed' in the sense that $(f,g): (N_1, H_1, \ell_1) \to (N_2, H_2, \ell_2)$ is sent to $(f,g): (N_1, H_1, \psi_{H_1}\ell_1) \to (N_2, H_2, \psi_{H_2}\ell_2)$, which may be seen to be a morphism in $\operatorname{Gl}(F_1)$ using the naturality of ψ .

We claim that Ψ^* is left adjoint to Ψ . To see this we consider the candidate counit $\varepsilon_{N,H,\ell} = (\varepsilon_{\ell}, \mathrm{id}_H)$, where ε_{ℓ} is defined as in the following pullback diagram.



We must show that given a morphism $(f,g): (N_1, H_1, \psi_{H_1}\ell_1) = \Psi^*(N_1, H_1, \ell_1) \to (N_2, H_2, \ell_2)$ there exists a unique morphism $(\widehat{f}, \widehat{g}): (N_1, H_1, \ell_1) \to \Psi(N_2, H_2, \ell_2) = (\overline{N_2}, H_2, \overline{\ell_2})$ such that $(\varepsilon_{\ell_2}, \operatorname{id}_{H_2}) \circ \Psi^*(\widehat{f}, \widehat{g}) = (f, g)$. We will construct this map using the following diagram.



Here the maps out of N_1 form a cone as we have $\psi_{H_2}F_2(g)\ell_1 = F_1(g)\psi_{H_1}\ell_1 = \ell_2 f$, where the first equality follows from naturality of ψ and the second from the fact that (f,g) is a morphism in $\text{Gl}(F_1)$.

By the universal property we have that $\overline{l_2}\hat{f} = F_2(g)\ell_1$, which means that (\hat{f}, g) is a morphism from (N_1, H_1, ℓ_1) to $(\overline{N_2}, H_2, \overline{\ell_2})$ in $\mathrm{Gl}(F_2)$. It is immediate from the diagram that $(\varepsilon_{\ell_2}, \mathrm{id}_{H_2}) \circ (\hat{f}, g) = (f, g)$ and it is also not hard to see that this is the unique such morphism. Thus, Ψ^* is indeed left adjoint to Ψ .

Finally, we must show that Ψ^* preserves finite limits. This follows immediately from the fact that finite limits in the glueing may be computed componentwise. \Box

Remark 4.14. Notice that Ψ^* is in fact a morphism of *non-split* extensions in the sense that it commutes with the kernel and cokernel maps up to isomorphism. However, it does not commute with the splittings unless Ψ is the identity.

4.3. Colimits in $Ext(\mathcal{H}, \mathcal{N})$

In [11] it was shown for frames H and N that there was something akin to a Baer sum of extensions in $\operatorname{Ext}(H, N)$. It is natural to ask if something analogous occurs in the category $\operatorname{Ext}(\mathcal{H}, \mathcal{N})$. Indeed, it is not hard to see via the equivalence with $\operatorname{Hom}(\mathcal{H}, \mathcal{N})^{\operatorname{op}}$ that $\operatorname{Ext}(\mathcal{H}, \mathcal{N})$ has all finite colimits. The following functor will help us compute these colimits.

Let $M: \operatorname{Ext}(\mathcal{H}, \mathcal{N})^{\operatorname{op}} \to \operatorname{Cat}/(\mathcal{N} \times \mathcal{H})$ be the functor sending extensions $\mathcal{N} \xrightarrow{K} \mathcal{G} \xleftarrow{E}_{E_*} \mathcal{H}$ to $!_{K^*E_*}^*: \operatorname{Gl}(K^*E_*) \to \mathcal{N} \times \mathcal{H}$ where $!_{K^*E_*}^*$ is left adjoint to the universal map $!_{K^*E_*}$ in $\operatorname{Ext}(\mathcal{H}, \mathcal{N})$ out of the 2-initial object $\mathcal{N} \xrightarrow{\pi_{1*}} \mathcal{N} \times \mathcal{H} \xleftarrow{\pi_{2*}} \mathcal{H}$. Explicitly, this adjoint sends (N, H, ℓ) to (N, H) and (f,g) to (f,g).

For the action on morphisms, let $(\Psi, \alpha, \beta, \gamma)$ be a morphism of extensions and let $\psi: K_2^* E_{2*} \to K_1^* E_{1*}$ be the corresponding natural transformation.



Then M maps Ψ to the morphism $\overline{\Psi}^*$: $\operatorname{Gl}(K_2^*E_{2*}) \to \operatorname{Gl}(K_1^*E_{1*})$, which sends (N, H, ℓ) to $(N, H, \psi_H \ell)$ and (f, g) to (f, g). It is immediate that the necessary diagram for this to be a morphism in the slice category commutes.

Proposition 4.15. The functor M above creates (and preserves) finite limits. Thus, finite colimits in $\text{Ext}(\mathcal{H}, \mathcal{N})$ can be computed from limits in $\text{Cat}/(\mathcal{N} \times \mathcal{H})$.

Proof. Let $D: \mathcal{J} \to \operatorname{Ext}(\mathcal{H}, \mathcal{N})$ be a diagram functor with finite domain \mathcal{J} . To compute the colimit of D we may compose D with $\Gamma_{\mathcal{H},\mathcal{N}}^{-1}$ and compute the limit in $\operatorname{Hom}(\mathcal{H},\mathcal{N})$. Let $R: \mathcal{H} \to \mathcal{N}$ be the resulting limit in $\operatorname{Hom}(\mathcal{H},\mathcal{N})$ and $(\varphi^i\colon R \to \Gamma_{\mathcal{H},\mathcal{N}}^{-1}D(i))_{i\in\mathcal{J}}$ be the corresponding projections. Then $\mathcal{N} \models \xrightarrow{\pi_{1*}} \operatorname{Gl}(R) \xrightarrow{\pi_{2*}} \mathcal{H}$ is the colimit of D, where the morphisms of the colimiting cone are given in the obvious way.

If we consider the diagram functor $MD: \mathcal{J} \to \operatorname{Cat}/(\mathcal{N} \times \mathcal{H})$, then we may again compute the limit with the assistance of the calculation in $\operatorname{Hom}(\mathcal{H}, \mathcal{N})$. We claim that $!_R^*: \operatorname{Gl}(R) \to \mathcal{N} \times \mathcal{H}$ is the required limit with the morphisms of the limiting cone given in the expected way — that is, if φ is a morphism of the limiting cone in $\operatorname{Hom}(\mathcal{H}, \mathcal{N})$ then $\Gamma_{\mathcal{H}, \mathcal{N}}(\varphi)^* = \overline{\Phi}^*$ is the associated morphism in $\operatorname{Cat}/(\mathcal{N} \times \mathcal{H})$.

We must demonstrate that this cone satisfies the universal property. Suppose we have some other cone $(\Xi_i: \mathcal{C} \to \operatorname{Gl}(K_i^* E_{i*}))_{i \in \mathcal{J}}$ and consider the following diagram in Cat where $\Psi = D(f)$ for some morphism $f: i \to j$ in \mathcal{J} .



Since each Ξ_i is a morphism in $\operatorname{Cat}/(\mathcal{N} \times \mathcal{H})$ we have that it commutes with the ! maps. This means that the Ξ maps all agree on the first two components. If we assume that $\Xi_k(C) = (N_C, H_C, \ell_C^k)$, then $\Xi_i = \overline{\Psi}^* \Xi_j$ gives $\ell_C^i = \psi_{H_C} \ell_C^j$. Now consider the following diagram in \mathcal{N} where we make use of the aforementioned limiting cone in $\operatorname{Hom}(\mathcal{H}, \mathcal{N})$.



Here we use the universal property of R componentwise at H_C to produce the map $\overline{\ell}_C$. This allows us to construct a map $S: \mathcal{C} \to \operatorname{Gl}(R)$ with $S(C) = (N_C, H_C, \overline{\ell}_C)$. As for morphisms, now note that each Ξ_k sends $f: C \to C'$ to the 'same' pair (f_1, f_2) and we define S to act on morphisms in the same way. The pair (f_1, f_2) can be seen to be a morphism in $\operatorname{Gl}(R)$ from S(C) to S(C') by considering the above diagram in the functor category and then using the naturality of $\overline{\ell}: \pi_1 \Xi_k \to R \pi_2 \Xi_k$. This morphism S is the desired map and is easily seen to be unique.

From the above it is clear that M preserves limits and that every limiting cone of MD is isomorphic to one of the form $(M\Gamma_{\mathcal{H},\mathcal{N}}(\varphi^i): M\Gamma_{\mathcal{H},\mathcal{N}}(R) \to MD(i))_{i \in \mathcal{J}}$, where $\varphi^i: R \to \Gamma_{\mathcal{H},\mathcal{N}}^{-1}D(i)$ is the limiting cone in $\operatorname{Hom}(\mathcal{H},\mathcal{N})$. For M to create limits, it remains to show that every cone of D which maps to a limiting cone of MD is isomorphic to one of the above form. This follows since M is conservative. \Box

Notice that the limit diagram was embedded into the slice category so that each Ξ in the proof would agree on the first two components. If the limit diagram is connected, this will happen automatically and so we obtain the following corollary.

Corollary 4.16. The functor $M: \operatorname{Ext}(\mathcal{H}, \mathcal{N}) \to \operatorname{Cat}$ sending extensions $\mathcal{N} \xrightarrow{K} \mathcal{G} \xleftarrow{E}_{E_*} \mathcal{H}$ to $\operatorname{Gl}(K^*E_*)$ and acting on morphisms as in Proposition 4.15 creates finite connected limits.

A disconnected (co)limit is the subject of the following example.

Example 4.17. Let us consider the coproduct of the extensions $\mathcal{N} \xrightarrow{K_1} \mathcal{G}_1 \xrightarrow{E_1} \mathcal{H}$ and

 $\mathcal{N} \xrightarrow{K_2} \mathcal{G}_2 \xleftarrow{E_2}{E_{2^*}} \mathcal{H}$. Since products in a slice category correspond to pullbacks, we may construct this coproduct using the following pullback in Cat.



If $!_{\mathcal{P}}$ is the composite morphism from \mathcal{P} to $\mathcal{N} \times \mathcal{H}$, then the coproduct extension may be recovered as $\mathcal{N} \xrightarrow{(\pi_1!_{\mathcal{P}})_*} \mathcal{P} \xrightarrow{\pi_2!_{\mathcal{P}}} \mathcal{H} \wedge$

$$P \xleftarrow{(\pi_2!_{\mathcal{P}})_*} H.$$

5. The Ext functor

Given that we have established that $\operatorname{Ext}(\mathcal{H}, \mathcal{N})$ is equivalent to $\operatorname{Hom}(\mathcal{H}, \mathcal{N})^{\operatorname{op}}$, it is natural to ask if Ext can be extended to a 2-bifunctor and if Ext and the $\operatorname{Hom}_{\operatorname{op}}$ will then be 2-naturally equivalent (where $\operatorname{Hom}_{\operatorname{op}} = \operatorname{Op} \circ \operatorname{Hom}^{\operatorname{co}}$ and Op is the opposite category 2-functor).

The answer is of course "yes", for if $T: \mathcal{H}' \to \mathcal{H}$ and $S: \mathcal{N} \to \mathcal{N}'$, all we need do is define $\operatorname{Ext}(T, S) = \Gamma_{\mathcal{H}',\mathcal{N}'} \circ \operatorname{Hom}_{\operatorname{op}}(T,S) \circ \Gamma_{\mathcal{H},\mathcal{N}}^{-1}$ (and similarly for natural transformations). However this is unsatisfactory, as there is already established behaviour for how an Ext functor ought to act on objects and morphisms. In this section, we show that the above definition conforms with the usual expectations of an Ext functor.

We consider each component of our Ext functor separately and begin by describing $\text{Ext}(-, \mathcal{N})$. In other contexts (for instance, see [3]) the extension functor can be obtained from a fibration. In the protomodular setting, we start from the 'fibration of points' sending split epimorphisms to their codomain. In the more general setting of \mathcal{S} -protomodularity (see [6]) we consider only a certain subclass of split epimorphisms. This suggests we consider a 2-fibration sending open subtopos adjunctions to the codomain of their inverse image functors.

A categorification of the Grothendieck construction (see [7]) gives that 2-fibrations correspond to 3-functors into 2Cat. Fortunately, aside from motivation, we will largely be able to avoid 3-functors for the same reasons that $\mathbb{E}xt(\mathcal{H},\mathcal{N})$ is essentially a 1-category (Corollary 4.5).

While the paradigmatic example of a fibration is the codomain fibration, which maps from the whole arrow category to the base category, the domain of the analogous 2-fibration is restricted to the category of fibrations. See [7] for more details on 2-fibrations.

The fibre 3-functor $\operatorname{Top}_{lex}^{co\,op} \to 2Cat$ corresponding to the 2-fibration Cod: $\operatorname{Fib}_{\operatorname{Top}_{lex}} \to \operatorname{Top}_{lex}$ can be described as follows (omitting the description of the coherence data for simplicity):

- On objects it sends a topos \mathcal{E} to the slice 2-category of finite-limit-preserving fibrations from toposes to \mathcal{E} .
- On 1-morphisms it sends a finite-limit-preserving functor $T: \mathcal{E}' \to \mathcal{E}$ to the 2-functor T^* corresponding to pulling back along T.
- On 2-morphisms it sends a natural transformation $\tau: T \to S$ to a 2-natural transformation τ^* from S^* to T^* . The component $\tau_E^*: S^*(E) \to T^*(E)$ indexed by the finite-limit-preserving fibration $E: \mathcal{D} \to \mathcal{E}$ can be constructed in three steps.

- First define the morphism of fibrations $(P_S, S): S^*(E) \to E$ given by the pullback projection.



- Then we may use the fibration property of E to lift the natural transformation $\tau S^{\star}(E)$ to a natural transformation into P_S . Explicitly, for an object $X \in \overline{\mathcal{D}}_S$, the morphism $\tau_{S^{\star}(E)(X)} \colon T(S^{\star}(E)(X)) \to S(S^{\star}(E)(X))$ can be lifted to a morphism in \mathcal{D} with codomain $P_S(X)$. These lifted morphisms as-

semble into a natural transformation $\overline{\tau}$ from a new functor $L: \overline{\mathcal{D}}_S \to \mathcal{D}$ to P_S . This functor sends a morphism $f: X \to Y$ to the morphism obtained by factoring $P_S(f)\overline{\tau}_X$ through $\overline{\tau}_Y$ as shown in the diagram below.



- Finally, consider the 2-pullback of E and T and note that the maps $L: \overline{\mathcal{D}}_S \to \mathcal{D}$ and $S^*(E): \overline{\mathcal{D}}_S \to \mathcal{E}'$ form a cone as shown below.



Thus, we may factor these through P_T and $T^*(E)$ respectively to obtain a functor from $\overline{\mathcal{D}}_S$ to $\overline{\mathcal{D}}_T$. This is the desired functor $\tau_E^* \colon S^*(E) \to T^*(E)$.

The coherent set of 2-isomorphisms for the 2-natural transformation can also be obtained by the cartesian property of the lifted maps and universal property of the 2-pullback.

We can then easily modify this to describe the fibre 3-functor for the 2-fibration of (open) points. Moreover, to obtain $\mathbb{E}xt(-,\mathcal{N})\colon \operatorname{Top}_{\operatorname{lex}}^{\operatorname{co\,op}} \to 2\operatorname{Cat}$ we restrict to inverse image functors of open subtoposes (equipped with right adjoint splittings) with fixed kernel object \mathcal{N} . The above discussion restricts easily to this case, since these functors are stable under pullback along finite-limit-preserving functors and the relevant morphisms can be shown to be morphisms of extensions. To see this we will use the following well-known result.

Proposition 5.1. Let C be a 2-category, $F: \mathcal{H} \to \mathcal{N}$ and $T: \mathcal{H}' \to \mathcal{H}$ 1-morphisms. Then the comma object Gl(FT) can be represented as a (strict) 2-pullback (where we draw the 2-cohernels of the extension *horizontally*).



Proof. Since Artin glueings are comma objects, this follows from the fact that a composition of comma square and a pullback square is a comma square. \Box

We will now describe the Ext 'functor' explicitly and at the same time demonstrate its relationship to Hom.

The 3-functor $\mathbb{E}xt(-,\mathcal{N})$ sends a topos \mathcal{H} to the 2-category of extensions $\mathbb{E}xt(\mathcal{H},\mathcal{N})$ as defined above. In fact, since $\mathbb{E}xt(\mathcal{H}, \mathcal{N})$ is equivalent to a 1-category, it turns out that $\mathbb{E}xt(-, \mathcal{N})$ factors through $Cat \hookrightarrow 2Cat$ and so can be treated as a 2-functor $\mathrm{Ext}(-,\mathcal{N})\colon \mathrm{Top}_{\mathrm{lex}}^{\mathrm{co\,op}}\to \mathrm{Cat}.$

The following computations will all be performed in the equivalent full subcategory of $Ext(\mathcal{H}, \mathcal{N})$ whose objects are of the form $\mathcal{N} \xrightarrow{\pi_{1*}} \operatorname{Gl}(F) \xrightarrow{\pi_2} \mathcal{H}$. The fact that this may be done coherently is due to the 2-categorical analogue of the result which says that if $F: \mathcal{B} \to \mathcal{C}$ is a functor and for each B we have an isomorphism $F(B) \cong C_B$, then there is a functor F' such that $F'(B) = C_B$ (and which acts on morphisms by conjugating the result of F by the appropriate isomorphisms).

We can use Proposition 5.1 to show how $\text{Ext}(-,\mathcal{N})$ acts on 1-morphisms. If $T: \mathcal{H}' \to \mathcal{H}$ is a finitelimit-preserving functor, then it is enough to describe how $\operatorname{Ext}(T, \mathcal{N}) \colon \operatorname{Ext}(\mathcal{H}, \mathcal{N}) \to \operatorname{Ext}(\mathcal{H}', \mathcal{N})$ acts on extensions in Artin-glueing form. It should 'take the 2-pullback along T' and hence it sends the object π^{FT}

$$\mathcal{N} \xrightarrow{\pi_{1*}^{F}} \operatorname{Gl}(F) \xrightarrow{\pi_{2}}^{F} \mathcal{H} \text{ to } \mathcal{N} \xrightarrow{\pi_{1*}^{F}} \operatorname{Gl}(FT) \xrightarrow{\pi_{2}}^{F} \mathcal{H}' \text{ as in the following diagram}$$



The functor $\operatorname{Ext}(T, \mathcal{N})$ acts on morphisms via the universal property of the 2-pullback. Let Ψ be a morphism of extensions corresponding to the natural transformation $\psi: F_2 \to F_1$ and consider the following diagram, noting that $T\pi_2^{F_1T} = \pi_2^{F_2}\Psi Q$ and hence we have a 2-cone. Also recall that Ψ can always be to chosen to correspond to pulling back along ψ .



We shall now show that we may take $\operatorname{Ext}(T, \mathcal{N})(\Psi)$ to be $(\Gamma_{\mathcal{H}', \mathcal{N}} \circ \operatorname{Hom}(T, \mathcal{N}) \circ \Gamma_{\mathcal{H}, \mathcal{N}}^{-1})(\Psi) = \Gamma_{\mathcal{H}', \mathcal{N}}(\psi T)$ — that is, the latter functor is given by the universal property of the 2-pullback. We must also supply two 2-morphisms corresponding to the left and top faces of the above cube. Both may be taken to be the identity. Note in fact that now each face of the cube commutes strictly.

We only need to check that $\pi_2^{F_2T}\Gamma_{\mathcal{H}',\mathcal{N}}(\psi T) = \pi_2^{F_1T}$ and that $Q'\Gamma_{\mathcal{H}',\mathcal{N}}(\psi T) = \Psi Q$. The former is immediate and the latter follows because $\Psi Q(N, H, \ell) = \Psi(N, T(H), \ell)$ is given by the following pullback.



This is readily seen to be the same pullback which determines $\Gamma_{\mathcal{H}',\mathcal{N}}(\psi T)(N,H,\ell)$.

Finally, we discuss how $\operatorname{Ext}(-, \mathcal{N})$ acts on 2-morphisms. We follow the construction outlined above for the codomain 2-fibration. Let $\tau: T \to T'$ be a natural transformation. Then we describe the natural transformation $\operatorname{Ext}(\tau, \mathcal{N}) : \operatorname{Ext}(T', \mathcal{N}) \to \operatorname{Ext}(T, \mathcal{N})$ componentwise.

Without loss of generality we may describe each component at extensions of the form $\Gamma_{\mathcal{H},\mathcal{N}}(F)$. Consider the following diagram.



As discussed in the case of the codomain fibration, we may define a functor L_{τ} as follows. We have that $P_{T'}(N, H, \ell) = (N, T'(H), \ell)$ lies above the codomain of $\tau_H \colon T(H) \to T'(H)$ with respect to the fibration π_2^F (see Proposition 2.5) and so by the universal property of the fibration we get a map $\overline{\tau}_{(N,H,\ell)} \colon (\overline{N}, T(H), \overline{\ell}) \to (N, T'(H), \ell)$ as given by the pair of vertical morphisms in the following pullback diagram.



These morphisms form a natural transformation $\overline{\tau} \colon L_{\tau} \to P_{T'}$, where L_{τ} is the functor which sends (N, H, ℓ) to $(\overline{N}, T(H), \overline{\ell})$. As above, this functor factors through P_T to give a functor $\operatorname{Ext}(\tau, \mathcal{N})_{\Gamma_{\mathcal{H},\mathcal{N}}(F)}$: $\operatorname{Gl}(FT') \to \operatorname{Gl}(FT)$, which sends (N, H, ℓ) to $(\overline{N}, H, \overline{\ell})$. It remains to show that this gives a morphism of split extensions, but it is clear from the above pullback diagram that this functor is the morphism of split extensions corresponding to τF (which itself is equal to $\operatorname{Hom}(\tau, \mathcal{N})_F$).

The morphisms $\operatorname{Ext}(\tau, \mathcal{N})_{\Gamma_{\mathcal{H},\mathcal{N}}(F)}$ define the desired natural transformation $\operatorname{Ext}(\tau, \mathcal{N})$. (Naturality follows from the interchange law or from the general theory of the codomain 2-fibration.)

The 2-functor $\operatorname{Ext}(-, \mathcal{N})$ composes strictly, and for the unitors observe that $\operatorname{Ext}(\operatorname{Id}_{\mathcal{H}}, \mathcal{N})$ is equal to $\Gamma_{\mathcal{H}, \mathcal{N}}\Gamma_{\mathcal{H}, \mathcal{N}}^{-1}$: $\operatorname{Ext}(\mathcal{H}, \mathcal{N}) \to \operatorname{Ext}(\mathcal{H}, \mathcal{N})$ so that we make take as our unitors the unit of the adjunction $\Gamma_{\mathcal{H}, \mathcal{N}}$. The necessary 2-functor axioms can then easily be seen to hold.

Remark 5.2. The above argument also proves that the 2-functor sending adjoint extensions with fixed kernel \mathcal{N} to their cokernels is a 2-fibration.

Theorem 5.3. The 2-functors $\operatorname{Ext}(-, \mathcal{N})$ and $\operatorname{Hom}_{\operatorname{op}}(-, \mathcal{N})$ are 2-naturally equivalent via $\Gamma^{\mathcal{N}}$: $\operatorname{Hom}_{\operatorname{op}}(-, \mathcal{N}) \to \operatorname{Ext}(-, \mathcal{N})$ defined as follows:

- (1) for each $\mathcal{H} \in \operatorname{Top}_{\operatorname{lex}}$ we have the equivalence $\Gamma^{\mathcal{N}}_{\mathcal{H}} = \Gamma_{\mathcal{H},\mathcal{N}} \colon \operatorname{Hom}(\mathcal{H},\mathcal{N})^{\operatorname{op}} \to \operatorname{Ext}(\mathcal{H},\mathcal{N}),$ (2) for each $T \colon \mathcal{H}' \to \mathcal{H}$ we have the identity $\Gamma_{\mathcal{H}',\mathcal{N}} \operatorname{Hom}(T,\mathcal{N})^{\operatorname{op}} = \operatorname{Ext}(T,\mathcal{N})\Gamma_{\mathcal{H},\mathcal{N}}.$

Proof. The equality in point (2) is clear by inspection of the definition of $\text{Ext}(T, \mathcal{N})$. The proof of the coherence conditions is easy. In particular, the first coherence condition is satisfied because each morphism of the diagram is the identity. Similarly, for the second condition again all morphisms are the identity (though marginally more work is required to show that the unitor at H whiskered with $\Gamma_{\mathcal{H},\mathcal{N}}$ is in fact the identity). \Box

We turn our attention to the functor $Ext(\mathcal{H}, -)$: Top_{lex}^{co} \rightarrow Cat. We could not find an elegant description of this in terms of a 2-fibration. However, we believe a reasonable definition can be given by dualising our arguments for $Ext(-, \mathcal{N})$.

Naturally, for an object \mathcal{N} we have that $\text{Ext}(\mathcal{H}, \mathcal{N})$ is just the category of adjoint split extensions.

For 1-morphisms, consider $S: \mathcal{N} \to \mathcal{N}'$. We would have $\operatorname{Ext}(\mathcal{H}, S)$ act on objects by sending $\mathcal{N} \xrightarrow{K} \mathcal{G} \xleftarrow{E}{\longleftarrow} \mathcal{H}$ to the extension resulting from a pushout of K along S and on morphisms via the universal property of the pushout.



To see that this is well-defined we prove the following result dual to Proposition 5.1.

Proposition 5.4. Let $F: \mathcal{H} \to \mathcal{N}$ and $S: \mathcal{N} \to \mathcal{N}'$ be finite-limit-preserving functors. Then Gl(SF) is given by the following 2-pushout.



Proof. By Proposition 4.7, we know that Gl(F) is a cocomma object in the 2-category Cat_{lex} . Thus it is a comma object in $\operatorname{Cat}_{\operatorname{lex}}^{\operatorname{op}}$. Applying Proposition 5.1 and then reversing the arrows gives the desired result. It is not hard to see that $P(N, H, \ell) = (S(N), H, S(\ell))$ and P(f, g) = (S(f), g). \Box

Thus, fixing particular pushouts we can describe $Ext(\mathcal{H}, S)$ concretely as sending an extension $\mathcal{N} \rightarrowtail \mathcal{G} \xleftarrow{E}{\longleftarrow} \mathcal{H} \text{ to } \mathcal{N} \rightarrowtail \operatorname{Gl}(SK^*E_*) \xleftarrow{\pi_2}{\longleftarrow} \mathcal{H}.$

As mentioned above, $\operatorname{Ext}(\mathcal{H}, S)$ should act on morphisms by the universal property of the pushout. Let $\Psi = (\Psi, \alpha_1, \beta_1, \gamma_1)$ be a morphism of extensions and consider the 2-cocone given by $P_2 \Psi \pi_{1*}^{F_1} \xrightarrow{P_2 \alpha_1^{-1}} P_2 \pi_{1*}^{F_2} \operatorname{Id}_{\mathcal{N}} = \pi_{1*}^{SF_2} \operatorname{Id}_{\mathcal{N}'} S$ as in the following pasting diagram (where we omit the 2-morphisms do avoid clutter).



Here the front, back and left faces commute strictly and the top face has associated invertible 2-morphism α_1^{-1} .

Let ψ be the natural transformation associated to Ψ . We will show that the map given by the universal property is $\Gamma_{\mathcal{H},\mathcal{N}'}(S\psi) = (\Gamma_{\mathcal{H},\mathcal{N}'}(S\psi), \alpha_2, \beta_2, \gamma_2)$. We define the associated 2-morphisms for the universal property of the 2-pushout as follows. For the bottom face we use $\alpha_2^{-1} \colon \Gamma_{\mathcal{H},\mathcal{N}'}(S\psi) \pi_{1*}^{SF_1} \to \pi_{1*}^{SF_2}$.

As for the right-hand face, we note that $P_2\Psi(N, H, \ell)$ is given by the following diagram.



On the other hand, we have that $\Gamma_{\mathcal{H},\mathcal{N}'}(S\psi)P_1(N,H,\ell)$ is given by a similar diagram as follows.



Since S preserves pullbacks, there is a natural family of isomorphisms $\mu_{(N,H,\ell)} = (\widehat{\mu}_{(N,H,\ell)}, \mathrm{id}_H)$ from $(S(\overline{N}), H, S(\overline{\ell}))$ to $(\overline{S(N)}, H, \overline{S(\ell)})$, which in particular satisfies that $\overline{S(\psi_H)}\widehat{\mu}_{(N,H,\ell)} = S(\overline{\psi_H})$. We take μ to be the 2-morphism associated to the right-hand face.

Now let us discuss the 2-cocone 2-morphism $P_2\alpha_1^{-1}$ in more detail. The component at N is given by $(S\widehat{\alpha}_{1,N}^{-1}, \mathrm{id})$ for $\widehat{\alpha}$ as defined Section 4.2. Notice that $S(\widehat{\alpha}_{1,N}^{-1})$ is given by the following diagram.



We see that $S(\widehat{\alpha}_1^{-1}) = S(\overline{\psi}_1)$. Similarly, we have $\overline{S(\psi_1)} = \widehat{\alpha}_2^{-1}S$. Thus, the equation $\overline{S(\psi_H)}\mu_N = S(\overline{\psi}_H)$ from above reduces in this case to $\widehat{\alpha}_{2,N}^{-1}S \circ \widehat{\mu}_{(N,1,!)} = S\widehat{\alpha}_{1,N}^{-1}$. Consequently, we have $\alpha_2^{-1}S \circ \mu \pi_{1*}^{F_1} = P_2\alpha_1^{-1}$, and hence $\Gamma_{\mathcal{H},\mathcal{N}}(S\psi)$ equipped with α_2^{-1} and (μ, id) is indeed the desired map from the universal property of the pushout.

In summary, we have $\operatorname{Ext}(\mathcal{H}, S)(\Psi) = \Gamma_{\mathcal{H}, \mathcal{N}'}(S\psi)$, which completes the description of $\operatorname{Ext}(\mathcal{H}, S)$.

Finally, let $\sigma: S \to S'$ be a natural transformation. We shall describe the natural transformation $\operatorname{Ext}(\mathcal{H}, \sigma): \operatorname{Ext}(\mathcal{H}, S') \to \operatorname{Ext}(\mathcal{H}, S)$ componentwise. We define $\operatorname{Ext}(\mathcal{H}, \sigma)_{\Gamma_{\mathcal{H},\mathcal{N}}(F)}$ through its left adjoint as follows.



Notice the resemblance of the above diagram to the one that arose when defining $\operatorname{Ext}(\tau, \mathcal{N})_{\Gamma_{\mathcal{H},\mathcal{N}}(F)}$. It has the same basic structure, but all 1-morphisms are pointing in the opposite direction.

First we define L_{σ} and $\overline{\sigma}$. Note that $P_{S'}(N, H, \ell) = (S'(N), H, S'(\ell))$ lies above the codomain of σ_N , with respect to the fibration $\pi_1^{S'F}$ (see Proposition 2.6). Thus by the universal property, we may lift σ_N to

a map $\overline{\sigma_N}$: $(S(N), H, S'(\ell)\sigma_N) \to (S'(N), H, S'(\ell))$. We may define a functor L_{σ} : $\operatorname{Gl}(F) \to \operatorname{Gl}(S'F)$ which sends objects (N, H, ℓ) to $(S(N), H, S'(\ell)\sigma_N)$ and which sends morphisms (f, g): $(N, H, \ell) \to (N', H', \ell')$ to (S(f), g). The pair (S(f), g) can be seen to be a morphism in $\operatorname{Gl}(S'F)$ by considering the following diagram.

$$\begin{array}{c|c} S(N) & \xrightarrow{\sigma_N} & S'(N) & \xrightarrow{S'(\ell)} & S'F(H) \\ \hline S(f) & & S'(f) & & & \\ & & & \\ & & \\ S(N') & \xrightarrow{\sigma_{N'}} & S'(N') & \xrightarrow{S'(\ell')} & S'F(H') \end{array}$$

The left-hand square commutes by naturality of σ and the right-hand square commutes since (f, g) is morphism in Gl(F). Now the $\overline{\sigma_N}$ arrange into a natural transformation $\overline{\sigma} \colon L_{\sigma} \to P_{S'}$.

If things are to behave dually, we should have L_{σ} factor through P_S . By the naturality of σ we have that the following diagram commutes.



Thus, observe that $L_{\sigma}(N, H, \ell) = (S(N), H, S'(\ell)\sigma_N) = (S(N), H, \sigma_{F(H)}S(\ell))$. This perspective allows us to factor L_{σ} as $L_{\sigma} = \Gamma_{\mathcal{H},\mathcal{N}'}(\sigma F)^* \circ P_S$. (Recall that the left adjoint to the functor $\Gamma_{\mathcal{H},\mathcal{N}'}(\sigma F)$ sends an object (N, H, ℓ) to $(N, H, \sigma_{F(H)}\ell)$.)

We might now hope to take $\operatorname{Ext}(\mathcal{H}, \sigma)_{F(H)}$ to be this resulting factor $\Gamma_{\mathcal{H},\mathcal{N}'}(\sigma F)^*$. However, this map goes in the 'wrong' direction. We can remedy this by taking the right adjoint and setting $\operatorname{Ext}(\mathcal{H}, \sigma)_{F(H)} = \Gamma_{\mathcal{H},\mathcal{N}'}(\sigma F)$.

In order to specify $\operatorname{Ext}(\mathcal{H}, -)$ completely, it only remains to discuss the compositors and unitors. As before we have that $\operatorname{Ext}(\mathcal{H}, -)$ composes strictly and we take the unitors to be the unit of the adjunction $\Gamma^*_{\mathcal{H},\mathcal{N}} \dashv \Gamma_{\mathcal{H},\mathcal{N}}$ in Theorem 4.12.

We can express the relationship between $Ext(\mathcal{H}, -)$ and $Hom(\mathcal{H}, -)$ as follows.

Theorem 5.5. $\operatorname{Ext}(\mathcal{H}, -)$ and $\operatorname{Hom}_{\operatorname{op}}(\mathcal{H}, -)$ are 2-naturally equivalent via $\Gamma^{\mathcal{H}} \colon \operatorname{Hom}_{\operatorname{op}}(\mathcal{H}, -) \to \operatorname{Ext}(\mathcal{H}, -)$ defined as follows

- (1) for each $\mathcal{N} \in \operatorname{Top}_{\operatorname{lex}}$ we have the equivalence $\Gamma_{\mathcal{N}}^{\mathcal{H}} = \Gamma_{\mathcal{H},\mathcal{N}} \colon \operatorname{Hom}(\mathcal{H},\mathcal{N})^{\operatorname{op}} \to \operatorname{Ext}(\mathcal{H},\mathcal{N}),$
- (2) for each $S: \mathcal{N} \to \mathcal{N}'$ we have the identity $\Gamma_{\mathcal{H},\mathcal{N}} \operatorname{Hom}(\mathcal{H},S)^{\operatorname{op}} = \operatorname{Ext}(\mathcal{H},S)\Gamma_{\mathcal{H},\mathcal{N}'}$.

Proof. Just as before, the equality in point (2) is clear by inspection of the definition of $\text{Ext}(\mathcal{H}, S)$ and the necessary coherence conditions hold, because each involved morphism is an identity. \Box

A bifunctor theorem for 2-functors was discussed in [10] and gives the precise conditions that allow two families of 2-functors $M_B: \mathcal{C} \to \mathcal{D}$ and $L_C: \mathcal{B} \to \mathcal{D}$ to be collated into a bifunctor $P: \mathcal{B} \times \mathcal{C} \to \mathcal{D}$ for which P(B, -) is isomorphic to M_B and P(-, C) is isomorphic to L_C . These conditions are that $L_C(B) = M_B(C)$ and that for each $f: B_1 \to B_2$ in \mathcal{B} and $g: C_1 \to C_2$ in \mathcal{C} there exists an invertible 2-morphism $\chi_{f,g}: L_{C_2}(g)M_{B_1}(f) \to M_{B_2}(f)L_{C_1}(g)$ satisfying certain coherence conditions reminiscent of those for distributive laws for monads. Such families together with χ are called a *distributive law of 2-functors*.

For the families $\operatorname{Ext}(-, \mathcal{N})$ and $\operatorname{Ext}(\mathcal{H}, -)$ the first condition is immediate. Moreover, it is not hard to see that if $T: \mathcal{H}' \to \mathcal{H}$ and $S: \mathcal{N} \to \mathcal{N}'$ that $\operatorname{Ext}(\mathcal{H}, S)\operatorname{Ext}(T, \mathcal{N}) = \operatorname{Ext}(T, \mathcal{N}')\operatorname{Ext}(\mathcal{H}', S)$ so that we might choose $\chi_{T,S}$ to be the identity. The coherence conditions are then immediate.

Thus, we may apply the results of [10] to arrive at the 2-functor (Ext, ω, κ) : Top_{lex}^{op} × Top_{lex} \rightarrow Cat defined below.

Definition 5.6. Let (Ext, ω, κ) : Top_{lex}^{co op} × Top_{lex}^{co} \rightarrow Cat be the 2-functor defined as follows.

- (1) $\operatorname{Ext}(\mathcal{H}, \mathcal{N})$ is the category of extensions of \mathcal{H} by \mathcal{N} ,
- (2) $\operatorname{Ext}(T, S) = \operatorname{Ext}(T, \mathcal{N}') \operatorname{Ext}(\mathcal{H}', S)$ for functors $S \colon \mathcal{N} \to \mathcal{N}'$ and $T \colon \mathcal{H}' \to \mathcal{H}$,
- (3) $\operatorname{Ext}(\tau, \sigma) = \operatorname{Ext}(\tau, \mathcal{N}') * \operatorname{Ext}(\mathcal{H}', \sigma)$ for 2-morphisms $\sigma \colon S \to S'$ and $\tau \colon T \to T'$,
- (4) ω is the identity,

(5) $\kappa_{\mathcal{H},\mathcal{N}}$ is given by the unit Φ of the adjunction $\Gamma^*_{\mathcal{H},\mathcal{N}} \dashv \Gamma_{\mathcal{H},\mathcal{N}}$ as defined in Theorems 3.17 and 4.12.

The 2-bifunctor $\operatorname{Hom}_{\operatorname{op}}$ can be recovered as the collation of the functors obtained by fixing one of its components, $\operatorname{Hom}_{\operatorname{op}}(\mathcal{H}, -)$ and $\operatorname{Hom}_{\operatorname{op}}(-, \mathcal{N})$. It is shown in [10] that 'morphisms between distributive laws' can also be collated to give 2-natural transformations between the corresponding bifunctors. The 2-natural equivalences in Theorems 5.3 and 5.5 can be collected into a 2-natural equivalence Γ : Ext \rightarrow Hom_{op} provided that $\Gamma_{\mathcal{H}}^{\mathcal{N}} = \Gamma_{\mathcal{N}}^{\mathcal{H}}$ and the Yang–Baxter equation holds. These conditions are immediate in our setting and so we obtain the following theorem.

Theorem 5.7. The 2-functors Ext and $\operatorname{Hom}_{\operatorname{op}}$ are 2-naturally equivalent via $\Gamma \colon \operatorname{Ext} \to \operatorname{Hom}_{\operatorname{op}}$ in which $\Gamma_{\mathcal{H},\mathcal{N}} = \Gamma_{\mathcal{N}}^{\mathcal{H}}$ and $\Gamma_{T,S}$ is the identity for all functors $S \colon \mathcal{N} \to \mathcal{N}'$ and $T \colon \mathcal{H}' \to \mathcal{H}$.

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